## Technical Document

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## Features

- Operating voltage: $\mathrm{f}_{\mathrm{SYS}}=4 \mathrm{MHz}: 2.2 \mathrm{~V} \sim 5.5 \mathrm{~V}$ $\mathrm{f}_{\mathrm{SYS}}=8 \mathrm{MHz}: 3.3 \mathrm{~V} \sim 5.5 \mathrm{~V}$
- 32 bidirectional I/O lines (PA, PB, PC, PD)
- Two external interrupt inputs
- Two 16-bit programmable timer/event counters with PFD (programmable frequency divider) function
- One 8-bit Remote Control Timer (RMT)
- Single channel serial interface
- VFD driver with $12 \times 16$ segments (12-segment \& 16 -grid to 20 -segment \& 8 -grid)
- $8 \mathrm{~K} \times 16 \times 2$ program memory
- $192 \times 8 \times 4$ data memory RAM
- Supports PFD for sound generation
- Real Time Clock (RTC), 32768Hz with quick start-up control bit


## General Description

The HT49RV9/HT49CV9 are 8-bit high performance single chip MCUs. Their single cycle instruction and 2-stage pipeline architecture make them suitable for high speed applications. As the devices include an VFD

- 8-bit prescaler for RTC
- Watchdog Timer
- Buzzer output
- On-chip crystal, RC and 32768 Hz crystal oscillator
- HALT function and wake-up feature reduce power consumption
- 16-level subroutine nesting
- 8-channel 10-bit resolution A/D converter
- 4-channel 8-bit PWM output shared with 4 I/O lines
- LVR function with enable/disable function
- Bit manipulation instruction
- 16-bit table read instruction
- Up to $0.5 \mu$ s instruction cycle with 8 MHz system clock
- 63 powerful instructions
- All instructions in 1 or 2 machine cycles
- 100-pin QFP package
driver they are suitable for use in products which require a front panel for their operation such as DVDs, VCDs, Mini-component audio systems, cassette decks, tuners, CD players, other home appliances, etc.


## Block Diagram



## Pin Assignment



Note: Each $V_{D D}\left(V_{S S}\right)$ pins must be connected to the power (ground) of the system.

## Pin Description

| Pin Name | I/O | Options | Description |
| :--- | :--- | :--- | :--- |
| PA0/BZ <br> PA1/BZ <br> PA2 <br> PA3/PFD <br> PA4~PA7 | I/O | Wake-up <br> Pull-high <br> Buzzer <br> PFD | Bidirectional 8-bit input/output port. Each bit can be configured as a <br> wake-up input by configuration option. Software instructions determine if <br> the pin is a CMOS output or Schmitt trigger input with or without pull-high <br> resistor (determined by pull-high options: bit option). Pins PA0, PA1 and <br> PA3 are pin-shared with BZ, BZ and PFD, respectively. |
|  |  | Bidirectional 8-bit input/output port. Software instructions determine if <br> the pin is a CMOS output or Schmitt trigger input with or without pull-high <br> resistor (determined by pull-high option: bit option) or A/D input. Once a <br> PB line is selected as an A/D input (by using software control), the I/O <br> function and pull-high resistor are disabled automatically. |  |
| PB0/AN0~ <br> PB7/AN7 | I/O | Pull-high |  |
| PC0~PC7 | I/O | Bidirectional 8-bit input/output port. Software instructions determine if <br> the pin is a CMOS output or Schmitt trigger input with or without pull-high <br> resistor (determined by pull-high option: bit option). |  |
| PD0/PWM0 <br> PD1/PWM1 <br> PD2/PWM2 <br> PD3/PWM3 | I/O | Pull-high | Bidirectional 4-bit input/output port. Software instructions determine if <br> the pin is a CMOS output or Schmitt trigger input with or without pull-high <br> resistor (determined by pull-high option: bit option). |


| Pin Name | I/O | Options | Description |
| :--- | :---: | :---: | :--- |
| PD4/INT0 <br> PD5/INT1 <br> PD6/TMR0 <br> PD7/TMR1 | I/O | Pull-high | Bidirectional 4-bit input/output port. Software instructions determine if <br> the pin is a CMOS output or Schmitt trigger input with or without pull-high <br> resistor (determined by pull-high option: bit option). Pins PD4~PD7 are <br> pin-shared with INT0, INT1, TMR0 \& TMR1, respectively (determined by <br> software control). |
| RMT | I | Pull-high | RMT with wake-up function (both rising and falling edge) and Schmitt <br> trigger input with or without pull-high resistor (determined by pull-high <br> option). |
| VSS | - | - | Negative power supply, ground |
| VEE | - | - | VFD negative power supply |
| SEG0~SEG11 | O | - | High-voltage segment output for VFD panel. |
| SEG12/GRID15~ <br> SEG19/GRID8 | O | - | High-voltage output for VFD panel. These pins are selectable for seg- <br> ment or grid output. |
| GRID0~GRID7 | O | - | High-voltage grid output for VFD panel. |$|$| I |
| :--- |
| SDI |

## Absolute Maximum Ratings

| Supply V | $\mathrm{V}_{\mathrm{SS}}-0.3 \mathrm{~V}$ to $\mathrm{V}_{\mathrm{SS}}+6.0 \mathrm{~V}$ | Storage Temperature ......................... $-50^{\circ} \mathrm{C}$ to $125^{\circ} \mathrm{C}$ |
| :---: | :---: | :---: |
| Input Voltage | $\mathrm{V}_{\text {SS }}-0.3 \mathrm{~V}$ to $\mathrm{V}_{\mathrm{DD}}+0.3 \mathrm{~V}$ | Operating Temperature ........................ $-40^{\circ} \mathrm{C}$ to $85^{\circ} \mathrm{C}$ |

Note: These are stress ratings only. Stresses exceeding the range specified under "Absolute Maximum Ratings" may cause substantial damage to the device. Functional operation of this device at other conditions beyond those listed in the specification is not implied and prolonged exposure to extreme conditions may affect device reliability.
D.C. Characteristics
$\mathrm{Ta}=25^{\circ} \mathrm{C}$

| Symbol | Parameter | Test Conditions |  | Min. | Typ. | Max. | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | V ${ }_{\text {D }}$ | Conditions |  |  |  |  |
| $V_{D D}$ | Operating Voltage | - | $\mathrm{f}_{\mathrm{SYS}}=4 \mathrm{MHz}$ | 2.2 | - | 5.5 | V |
|  |  | - | $\mathrm{f}_{\mathrm{SYS}}=8 \mathrm{MHz}$ | 3.3 | - | 5.5 | V |
| VEE | VFD Supply Voltage | - | - | 0 | - | $\mathrm{V}_{\mathrm{DD}}-30$ | V |
| $\mathrm{l}_{\mathrm{DD} 1}$ | Operating Current (Crystal OSC) | 3 V | No load, ADC off VFD off, fsYS $=4 \mathrm{MHz}$ | - | 2 | 3 | mA |
|  |  | 5 V |  | - | 5 | 8 | mA |
| $\mathrm{I}_{\mathrm{DD} 2}$ | Operating Current (RC OSC) | 3 V | No load, ADC off VFD off, $\mathrm{f}_{\mathrm{SYS}}=4 \mathrm{MHz}$ | - | 1.8 | 2.7 | mA |
|  |  | 5 V |  | - | 4.6 | 7.5 | mA |
| $\mathrm{I}_{\mathrm{DD} 3}$ | Operating Current ( $f_{S Y S}=32768 \mathrm{~Hz}$ ) | 3V | No load, ADC off VFD off | - | 1.2 | 2 | mA |
|  |  | 5 V |  | - | 4 | 7 | mA |
| $\mathrm{I}_{\mathrm{DD} 4}$ | Operating Current (Crystal OSC) | 3 V | No load, ADC off VFD on, , $\mathrm{f}_{\mathrm{SYS}}=4 \mathrm{MHz}$ | - | 4 | 6 | mA |
|  |  | 5 V |  | - | 7 | 15 | mA |
| ISTB1 | Standby Current ( ${ }^{*} \mathrm{f}_{\mathrm{S}}=\mathrm{T} 1$ ) | 3V | No load, system HALT VFD off at HALT | - | - | 1 | $\mu \mathrm{A}$ |
|  |  | 5 V |  | - | - | 2 | $\mu \mathrm{A}$ |
| ISTB2 | Standby Current$\left({ }^{*} f_{\mathrm{S}}=32768 \mathrm{~Hz} \text { OSC }\right)$ | 3 V | No load, system HALT VFD off at HALT | - | 4 | 10 | $\mu \mathrm{A}$ |
|  |  | 5 V |  | - | 14 | 20 | $\mu \mathrm{A}$ |
| $V_{\text {IL } 1}$ | Input Low Voltage for I/O Ports, TMR and INT | - | - | 0 | - | $0.2 \mathrm{~V}_{\mathrm{DD}}$ | V |
| $\mathrm{V}_{\mathrm{IH} 1}$ | Input High Voltage for I/O Ports, TMR and INT | - | - | 0.8 V DD | - | $V_{D D}$ | V |
| $\mathrm{V}_{\text {IL2 }}$ | Input Low Voltage ( $\overline{\mathrm{RES}}$ ) | - | - | 0 | - | $0.4 \mathrm{~V}_{\mathrm{DD}}$ | V |
| $\mathrm{V}_{\mathrm{IH} 2}$ | Input High Voltage ( $\overline{\mathrm{RES}}$ ) | - | - | $0.9 \mathrm{~V}_{\mathrm{DD}}$ | - | $V_{D D}$ | V |
| $\mathrm{V}_{\text {LVR }}$ | Low Voltage Reset Voltage | - | LVR voltage 3.0V option | 2.7 | 3.0 | 3.3 | V |
|  |  |  | LVR voltage 3.8 V option | 3.5 | 3.8 | 4.0 | V |
| Iol | I/O Port Segment Logic Output Sink Current | 3V | $\mathrm{V}_{\mathrm{OL}}=0.1 \mathrm{~V}_{\mathrm{DD}}$ | 6 | 12 | - | mA |
|  |  | 5 V |  | 10 | 25 | - | mA |
| IOH | I/O Port Segment Logic Output Source Current | 3V | $\mathrm{V}_{\mathrm{OH}}=0.9 \mathrm{~V}_{\mathrm{DD}}$ | -2 | -4 | - | mA |
|  |  | 5 V |  | -5 | -8 | - | mA |
| $\mathrm{IOH2}$ | Segment/Grid Source Current | 5 V | $\mathrm{V}_{\mathrm{OH}}=\mathrm{V}_{\mathrm{DD}}-2 \mathrm{~V}$ | -15 | - | - | mA |
| $\mathrm{R}_{\mathrm{PH}}$ | Pull-high Resistance of I/O Ports and $\overline{\mathrm{NTTO}}, \overline{\mathrm{INT}}, \mathrm{RMT}$ | 3 V | - | 40 | 60 | 80 | k $\Omega$ |
|  |  | 5 V | - | 10 | 30 | 50 | $\mathrm{k} \Omega$ |
| $\mathrm{V}_{\mathrm{AD}}$ | A/D Input Voltage | - | - | 0 | - | $V_{D D}$ | V |
| $E_{\text {AD }}$ | A/D Conversion Integral Nonlinearity Error | - | - | - | $\pm 0.5$ | $\pm 1$ | LSB |
| $\mathrm{I}_{\text {ADC }}$ | Additional Power Consumption if $A / D$ Converter is Used | 3V | - | - | 1 | 2 | mA |
|  |  | 5 V |  | - | 2 | 4 | mA |

Note: "*fs" Refer to WDT clock option

## A.C. Characteristics

$\mathrm{Ta}=25^{\circ} \mathrm{C}$

| Symbol | Parameter | Test Conditions |  | Min. | Typ. | Max. | Unit |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\mathrm{V}_{\mathrm{DD}}$ | Conditions |  |  |  |  |
| $\mathrm{f}_{\text {SYS }} 1$ | System Clock | - | 2.2V~5.5V | 400 | - | 4000 | kHz |
|  |  | - | $3.3 \mathrm{~V} \sim 5.5 \mathrm{~V}$ | 400 | - | 8000 | kHz |
| $\mathrm{f}_{\text {SYS } 2}$ | System Clock (32768Hz Crystal OSC) | - | 2.2V $\sim 5.5 \mathrm{~V}$ | - | 32768 | - | Hz |
| $\mathrm{f}_{\text {RTCOSC }}$ | RTC Frequency | - | - | - | 32768 | - | Hz |
| $\mathrm{f}_{\text {TIMER }}$ | Timer I/P Frequency (TMR0/TMR1) | - | 2.2V~5.5V | 0 | - | 4000 | kHz |
|  |  | - | $3.3 \mathrm{~V} \sim 5.5 \mathrm{~V}$ | 0 | - | 8000 | kHz |
| twdTosc | Watchdog Oscillator Period | 3V | - | 45 | 90 | 180 | $\mu \mathrm{s}$ |
|  |  | 5 V | - | 32 | 65 | 130 | $\mu \mathrm{s}$ |
| $t_{\text {RES }}$ | External Reset Low Pulse Width | - | - | 1 | - | - | $\mu \mathrm{S}$ |
| $\mathrm{t}_{\text {SST }}$ | System Start-up Timer Period | - | Power-up or wake-up from HALT | - | 1024 | - | $\mathrm{t}_{\text {SYS }}$ |
| $\mathrm{t}_{\text {INT }}$ | Interrupt Pulse Width | - | - | 1 | - | - | $\mu \mathrm{s}$ |
| $t_{\text {AD }}$ | A/D Clock Period | - | - | 1 | - | - | $\mu \mathrm{s}$ |
| $\mathrm{t}_{\text {ADC }}$ | A/D Conversion Time | - | - | - | 76 | - | $\mathrm{t}_{\text {AD }}$ |
| $t_{\text {ADCS }}$ | A/D Sampling Time | - | - | - | 32 | - | $t_{\text {AD }}$ |

Note: $\quad t_{S Y S}=1 / \mathrm{f}_{S Y S}$

## Functional Description

## Execution Flow

The system clock is derived from either a crystal or an RC oscillator or a 32768 Hz crystal oscillator. It is internally divided into four non-overlapping clocks. One instruction cycle consists of four system clock cycles.

Instruction fetching and execution are pipelined in such a way that a fetch takes one instruction cycle while decoding and execution takes the next instruction cycle. The pipelining scheme ensures that instructions are effectively executed in one cycle. Exceptions to this are instructions that change the contents of the program counter, such as subroutine calls or jumps, in which case, two cycles are required to complete the instruction.

## Program Counter - PC

The 13-bit program counter (PC) controls the sequence in which the instructions stored in the program ROM are
executed and its contents specify a maximum of 8192 addresses.

After accessing a program memory word to fetch an instruction code, the contents of the program counter are incremented by one. The program counter then points to the memory word containing the next instruction code.

When executing a jump instruction, conditional skip execution, loading PCL register, subroutine call, initial reset, internal interrupt, external interrupt or return from subroutine, etc., the microcontroller manages program control by loading the address corresponding to each instruction.

For conditional skip instructions, once the condition has been met, the next instruction, which has already been fetched during the current instruction execution, is discarded and a dummy cycle replaces it while the proper instruction is obtained. Otherwise proceed with the next instruction.


Execution Flow

| Mode | Program Counter |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | *12 | *11 | *10 | *9 | *8 | *7 | *6 | *5 | *4 | *3 | *2 | *1 | *0 |
| Initial Reset | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| External Interrupt 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| External Interrupt 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| Timer/Event Counter 0 Overflow | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 |
| Timer/Event Counter 1 Overflow | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| Serial Interface Interrupt | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | 0 | 0 |
| Multi-function Interrupt | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 |
| Skip | Program Counter+2 |  |  |  |  |  |  |  |  |  |  |  |  |
| Loading PCL | *12 | *11 | *10 | *9 | *8 | @7 | @6 | @5 | @4 | @3 | @2 | @1 | @0 |
| Jump, Call Branch | \#12 | \#11 | \#10 | \#9 | \#8 | \#7 | \#6 | \#5 | \#4 | \#3 | \#2 | \#1 | \#0 |
| Return from Subroutine | S12 | S11 | S10 | S9 | S8 | S7 | S6 | S5 | S4 | S3 | S2 | S1 | S0 |

## Program Counter

Note: ${ }^{* 12 ~ *} 0$ : Program counter bits \#12~\#0: Instruction code bits

S12~S0: Stack register bits
@7~@0: PCL bits

The lower byte of the program counter (PCL) is available for program control and is a readable and writeable register ( 06 H ). Moving data into the PCL performs a short jump. The destination will be within 256 locations.

When a control transfer takes place, an additional dummy cycle is required.

## Program Memory

The program memory is used to store the program instructions which are to be executed. It also contains data, table, and interrupt entries, and is organized into $16384 \times 16$ bits format. The program counter is composed of 13 bits, so it can directly access the whole program memory without changing banks.

The ROM bank 0 ranges from 0000 H to 1 FFFH ( $\mathrm{BP}=000 \mathrm{XXXXXB}$ ). The ROM bank1 ranges from 2000 H to $3 F F F H$ ( $B P=001 X X X X X B$ ).
Certain locations in the ROM are reserved for special usage:

- Location 000 H

This area is reserved for use by the chip reset for program initialization. After a chip reset is initiated, the program will jump to this location and begin execution.


Program Memory

- Location 004H

This area is reserved for the external interrupt service program. If the $\overline{\mathrm{INTO}}$ input pin is activated, and the interrupt is enabled, and the stack is not full, the program will jump to this location and begin execution.

- Location 008H

This area is reserved for the external interrupt service program. If the $\overline{\mathrm{INT}}$ input pin is activated, and the interrupt is enabled, and the stack is not full, the program will jump to this location and begin execution.

- Location 00CH

This area is reserved for the Timer/Event Counter 0 interrupt service program. If a timer interrupt results from a Timer/Event Counter 0 overflow, and if the interrupt is enabled and the stack is not full, the program will jump to this location and begin execution.

- Location 010H

This area is reserved for the Timer/Event Counter 1 interrupt service program. If a timer interrupt results from a Timer/Event Counter 1 overflow, and if the interrupt is enabled and the stack is not full, the program will jump to this location and begin execution.

- Location 014H

This area is reserved for the serial interface interrupt service program. If 8 bits of data have been received or transmitted successfully from the serial interface, and the interrupt is enabled, and the stack is not full, the program will jump to this location and begin execution.

- Location 018H

This area is reserved for the multi-function interrupt. If a real time clock interrupt occurs, or if a rising edge is detected from the RMT input pin, or if a falling edge is detected from the RMT input pin, or if the RMT overflow and the related interrupts are enabled, and the stack is not full, the program will jump to this location and begin execution.

- Table location

Any location within the program memory can be used as a look-up table where programmers can store fixed data. The instructions "TABRDC [m]" (the current page, 1 page=256 words) and "TABRDL [m]" (the last page) transfer the contents of the lower-order byte to the specified data memory, and the contents of the higher-order byte to TBLH (Table Higher-order byte register) $(08 \mathrm{H})$. Only the destination of the lower-order byte in the table is well-defined, the other bits of the ta-

| Instruction(s) | Table Location |  |  |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | *12 | *11 | *10 | *9 | *8 | *7 | *6 | *5 | *4 | *3 | *2 | *1 | *0 |
| TABRDC [m] | P12 | P11 | P10 | P9 | P8 | @7 | @6 | @ 5 | @4 | @ 3 | @2 | @1 | @0 |
| TABRDL [m] | 1 | 1 | 1 | 1 | 1 | @7 | @6 | @5 | @4 | @3 | @ 2 | @1 | @0 |

Table Location
Note: *12~*0: Table location bits
@7~@0: Table pointer bits
ble word are all transferred to the lower portion of TBLH. The TBLH is a read only register and the table pointer (TBLP) is a read/write register ( 07 H ), which indicates the table location. Before accessing the table, the location must be placed in the TBLP. All table related instructions require two cycles to complete the operation. These areas may function as normal program memory depending upon user's requirements.

## Stack Register - STACK

The stack register is a special part of the memory used to save the contents of the Program Counter. The stack is organized into 16 levels and is neither part of the data nor part of the program, and is neither readable nor writeable. Its activated level is indexed by a stack pointer (SP) and is neither readable nor writeable. At the start of a subroutine call or an interrupt acknowledgment, the contents of the Program Counter is pushed onto the stack. At the end of the subroutine or interrupt routine, signaled by a return instruction (RET or RETI), the contents of the Program Counter is restored to its previous value from the stack. After a chip reset, the SP will point to the top of the stack.

If the stack is full and a non-masked interrupt takes place, the interrupt request flag is recorded but the acknowledgment is still inhibited. Once the SP is decremented (by RET or RETI), the interrupt is serviced. This feature prevents stack overflow, allowing the programmer to use the structure easily. Likewise, if the stack is full, and a "CALL" is subsequently executed, a stack overflow occurs and the first entry is lost (only the most recent sixteen return addresses are stored).

## Data Memory - RAM

The data memory (RAM) has a capacity of $810 \times 8$ bits, and is divided into two functional groups, namely; special function registers ( $42 \times 8$ bit) and general purpose data memory (Bank 0, 2~4: each RAM bank contains $192 \times 8$ bits) most of which are readable/writeable, but some are read only. The special function registers are overlapped in any banks.

The special function registers consist of an Indirect addressing register $0(00 \mathrm{H})$, a Memory pointer register 0 (MP0;01H), an Indirect addressing register 1 (02H), a Memory pointer register 1 (MP1;03H), a Bank pointer (BP;04H), an Accumulator (ACC;05H), a Program counter lower-order byte register (PCL;06H), a Table pointer (TBLP;07H), a Table higher-order byte register (TBLH;08H), a Real time clock control register (RTCC;09H), a Status register (STATUS;0AH), an Interrupt control register 0 (INTC0;0BH), a Timer/Event Counter 0 (TMROH:OCH; TMROL:ODH), a Timer/Event Counter 0 control register (TMROC;0EH), a Timer/Event Counter 1 (TMR1H:0FH;TMR1L:10H), a Timer/Event Counter 1 control register (TMR1C; 11H), Interrupt control register 1 (INTC1;1EH), Serial bus control register


RAM Mapping
(SBCR;1FH), Serial bus data register (SBDR; 20H), Remote Timer control register (RMTC;21H), Multi-function interrupt status register 3 (MFIS;29H), PWM data register (PWM0;1AH, PWM1;1BH, PWM2;1CH, PWM3; 1DH), the A/D result lower-order byte register (ADRL;24H), the A/D result higher-order byte register (ADRH;25H), the A/D control register (ADCR;26H), the A/D clock setting register (ACSR;27H), VFD control register (VFDC;28H), I/O registers (PA;12H, PB;14H, PC;16H, PD;18H) and I/O control registers (PAC;13H, PBC; 15H, PCC; 17H and PDC;19H).

The remaining space before 40 H is reserved for future expanded usage and reading these locations will return the result " 00 H ". The space before 40 H overlaps in each bank. The general-purpose data memory, addressed from 40H to FFH (Bank0; BP=0 or Bank2; $\mathrm{BP}=2$ or Bank3; $\mathrm{BP}=3$ or Bank4; $\mathrm{BP}=4$ ), is used for data and control information under instruction commands.

All data memory areas can handle arithmetic, logic, increment, decrement and rotate operations directly. Except for some dedicated bits, each bit in the data memory can be set and reset by "SET [m].i" and "CLR [m].i". They are also indirectly accessible through memory pointer registers (MP0;01H/MP1;03H).

## Indirect Addressing Register

Location 00 H and 02 H are indirect addressing registers that are not physically implemented. Any read/write operation of $[00 \mathrm{H}]$ and $[02 \mathrm{H}]$ accesses the RAM pointed to by MP0 (01H) and MP1 (03H), respectively. Reading location 00 H or 02 H indirectly returns the result 00 H . Writing indirectly leads to no operation.

The function of data movement between two indirect addressing registers is not supported. The memory pointer registers, MP0 and MP1, are both 8-bit registers used to access the RAM by combining corresponding indirect addressing registers. MP0 can only be applied to data memory, while MP1 can be applied to data memory and VFD display memory.

## Accumulator - ACC

The accumulator (ACC) is closely related with operations carried out by the ALU. It is mapped to location 05 H of the RAM and is capable of operating with immediate data. The data movement between two data memory locations must pass through the ACC.

## Arithmetic and Logic Unit - ALU

This circuit performs 8-bit arithmetic and logic operations and provides the following functions:

- Arithmetic operations (ADD, ADC, SUB, SBC, DAA)
- Logic operations (AND, OR, XOR, CPL)
- Rotation (RL, RR, RLC, RRC)
- Increment and Decrement (INC, DEC)
- Branch decision (SZ, SNZ, SIZ, SDZ etc.)

The ALU not only saves the results of a data operation but also changes the status register.

## Status Register - STATUS

The status register ( OAH ) is 8 bits wide and contains, a carry flag (C), an auxiliary carry flag (AC), a zero flag (Z), an overflow flag (OV), a power down flag (PDF), and a watchdog time-out flag (TO). It also records the status information and controls the operation sequence.

Except for the TO and PDF flags, bits in the status register can be altered by instructions similar to other registers. Data written into the status register does not alter the TO or PDF flags. Operations related to the status register, however, may yield different results from those intended. The TO and PDF flags can only be changed by a Watchdog Timer overflow, chip power-up, or clearing the Watchdog Timer and executing a "HALT" instruction. The Z, OV, AC, and C flags reflect the status of the latest operations.

On entering an interrupt sequence or executing a subroutine call, the status register will not be automatically pushed onto the stack. If the contents of the status is important, and if the subroutine is likely to corrupt the status register, the programmer should take precautions and save it properly.

| Bit No. | Label | Function |
| :---: | :---: | :---: |
| 0 | C | $C$ is set if an operation results in a carry during an addition operation or if a borrow does not take place during a subtraction operation; otherwise C is cleared. C is also affected by a rotate through carry instruction. |
| 1 | AC | AC is set if an operation results in a carry out of the low nibbles in addition or no borrow from the high nibble into the low nibble in subtraction, otherwise AC is cleared. |
| 2 | Z | Z is set if the result of an arithmetic or logic operation is zero, otherwise Z is cleared. |
| 3 | OV | OV is set if an operation results in a carry into the highest-order bit but not a carry out of the highest-order bit, or vice versa, otherwise OV is cleared. |
| 4 | PDF | PDF is cleared by either a system power-up or executing the "CLR WDT" instruction. PDF is set by executing the "HALT" instruction. |
| 5 | TO | TO is cleared by a system power-up or executing the "CLR WDT" or "HALT" instruction. TO is set by a WDT time-out. |
| 6, 7 | - | Unused bit, read as "0" |

Status Register

## Interrupts

The devices provides two external interrupts, two internal timer/event counter interrupts, three remote control timer interrupts, an internal real time clock interrupt and serial interface interrupt. The interrupt control register 0 (INTCO;0BH) and interrupt control register 1 (INTC1;1EH) both contain the interrupt control bits that are used to set the enable/disable status and interrupt request flags.

Once an interrupt subroutine is serviced, all other interrupts are blocked (by clearing the EMI bit). This scheme may prevent any further interrupt nesting. Other interrupt requests may take place during this interval, but only the interrupt request flag will be recorded. If another interrupt requires servicing while the program is in the interrupt service routine, the programmer should set the EMI bit and the corresponding bit of the INTC0 or INTC1 to allow interrupt nesting. Once the stack is full, the interrupt request will not be acknowledged, even if the related interrupt is enabled, until the SP is decremented. If immediate service is desired, the stack should be prevented from becoming full.

All these interrupts can support a wake-up function. As an interrupt is serviced, a control transfer occurs by pushing the contents of the Program Counter onto the stack followed by a branch to a subroutine at the specified location in the ROM. Only the contents of the Program Counter is pushed onto the stack. If the contents of the register or of the status register (STATUS) is altered by the interrupt service program which corrupts the desired control sequence, the contents should be saved in advance.

External interrupts are triggered by an edge transition of INT0 or INT1 (configuration option: high to low, low to high, low to high or high to low), and the related interrupt request flag (EIF0; bit 4 of INTC0, EIF1; bit 5 of INTC0) is set as well. After the interrupt is enabled, the stack is not full, and the external interrupt is active, a subroutine call to location 04 H or 08 H occurs. The interrupt request flag (EIF0 or EIF1) and EMI bits are all cleared to disable other maskable interrupts.

The internal Timer/Event Counter 0 interrupt is initialized by setting the Timer/Event Counter 0 interrupt request flag (TOF; bit 6 of INTCO), which is normally caused by a timer overflow. After the interrupt is enabled, and the stack is not full, and the T0F bit is set, a subroutine call to location 0 CH occurs. The related interrupt request flag (TOF) is reset, and the EMI bit is cleared to disable other maskable interrupts. Timer/Event Counter 1 is operated in the same manner but its related interrupt request flag is T1F (bit 4 of INTC1) and its subroutine call location is 10 H .

The Serial Interface interrupt is initialized by setting the interrupt request flag (TDRF; bit 5 of INTC1), that is caused by completely receiving/transferring 8 bits of
data from a serial interface. After the interrupt is enabled, and the stack is not full, and the TRF bit is set, a subroutine call to location 14 H occurs. The related interrupt request flag (TDRF) is reset and the EMI bit is cleared to disable further maskable interrupts.

The multi-function interrupt is initialized by setting the interrupt request flag (MFF; bit 6 of INTC1), that is caused by a regular real time clock signal, caused by a rising edge of RMT, or caused by a falling edge of RMT or caused by an RMT overflow. After the interrupt is enabled, the stack is not full, and the MFF bit is set, a subroutine call to location 18 H occurs. The related interrupt request flag (MFF) is reset and the EMI bit is cleared to disable further maskable interrupts.
During the execution of an interrupt subroutine, other maskable interrupt acknowledgments are all held until the "RETI" instruction is executed or the EMI bit and the related interrupt control bit are both set to 1 (if the stack is not full). To return from the interrupt subroutine, "RET" or "RETI" may be invoked. RETI sets the EMI bit and enables an interrupt service, but RET does not.

Interrupts occurring in the interval between the rising edges of two consecutive T2 pulses are serviced on the latter of the two T2 pulses if the corresponding interrupts are enabled. In the case of simultaneous requests, the following table shows the priority that is applied. These can be masked by resetting the EMI bit.

| Interrupt Source | Priority | Vector |
| :--- | :---: | :---: |
| External interrupt 0 | 1 | 04 H |
| External interrupt 1 | 2 | 08 H |
| Timer/Event Counter 0 overflow | 3 | 0 CH |
| Timer/Event Counter 1 overflow | 4 | 10 H |
| Serial interface interrupt | 5 | 14 H |
| Multi-function interrupt | 6 | 18 H |

The RMT overflow interrupt flag (RMTVF; bit 0 of MFIS), real time clock interrupt flag (RTF; bit 1 of MFIS), the RMT rising edge interrupt flag (RMTOF; bit 2 of MFIS) and the RMT falling edge interrupt flag (RMT1F; bit 3 of MFIS) indicate that a related interrupt has occurred. After reading these flags, these flags will not be cleared automatically, they should be cleared by the user.

The serial interface interrupt is indicated by the interrupt flag (TDRF; bit 5 of INTC1), that is caused by receiving or transferring a complete 8-bit data transfer between the HT49RV9/ HT49CV9 and an external device. After the interrupt is enabled (by setting ESBI; bit 1 of INTC1), and the stack is not full, a subroutine call to location 14 H occurs.
The Timer/Event Counter 0 interrupt request flag (TOF), external interrupt 1 request flag (EIF1), external interrupt 0 request flag (EIF0), enable Timer/Event Counter0 interrupt bit (ETOI), enable external interrupt 1 bit (EEI1),

| Bit No. | Label | Function |
| :---: | :---: | :--- |
| 0 | EMI | Controls the master or global interrupt (1=enable; 0=disable) |
| 1 | EEIO | Controls the external interrupt 0 (1=enable; $0=$ disable) |
| 2 | EEI1 | Controls the external interrupt 1 (1=enable; 0=disable) |
| 3 | ETOI | Controls the Timer/Event Counter 0 interrupt (1=enable; 0=disable) |
| 4 | EIF0 | External interrupt 0 request flag (1=active; $0=$ inactive) |
| 5 | EIF1 | External interrupt 1 request flag (1=active; 0=inactive) |
| 6 | T0F | Internal Timer/Event Counter 0 request flag (1=active; $0=$ inactive) |
| 7 | - | Unused bit, read as "0" |

INTCO (OBH) Register

| Bit No. | Label | Function |
| :---: | :---: | :--- |
| 0 | ET1I | Controls the Timer/Event Counter 1 interrupt (1=enable; 0=disable) |
| 1 | ESBI | Controls the Serial Interface interrupt (1=enable; 0:disable) |
| 2 | EMFI | Controls the multi-function interrupt (1=enable; 0:disable) |
| 3,7 | - | Unused bit, read as "0" |
| 4 | T1F | Internal Timer/Event Counter 1 request flag (1=active; 0=inactive) |
| 5 | TDRF | Serial bus data transferred or data received interrupt request flag (1=active; 0=inactive) |
| 6 | MFF | Multi-function interrupt request flag (1=active; 0=inactive) |

INTC1 (1EH) Register

| Bit No. | Label | Function |
| :---: | :---: | :--- |
| 0 | - | Unused bit, read as " $0 " \prime$ |

RMTC (21H) Register

| Bit No. | Label | Function |
| :---: | :---: | :--- |
| 0 | RMTVF | Remote control timer overflow interrupt flag (1=indicates that an overflow has occurred; 0=indi- <br> cate that an overflow has not occurred) |
| 1 | RTF | Real time clock interrupt flag (1=indicates that an RTC interrupt has occurred; 0=indicates that <br> an RTC interrupt has not occurred) |
| 2 | RMT0F | Remote control timer rising edge interrupt flag (1=indicates that a rising edge interrupt has oc- <br> curred; 0=indicates that a rising edge interrupt has not occurred) |
| 3 | RMT1F | Remote control timer falling edge interrupt flag (1=indicates that a falling edge interrupt has oc- <br> curred; 0=indicates that a falling edge interrupt has not occurred) |
| 4 | ERTI | Controls the real time clock interrupt (1=enable; 0=disable) |
| $5 \sim 7$ | - | Unused bit, read as "0" |

MFIS (29H) Register
enable external interrupt 0 bit (EEIO), and enable master interrupt bit (EMI) constitute the Interrupt Control register 0 (INTCO) which is located at OBH in the RAM.
The multi-function interrupt request flag (MFF), serial interface interrupt request flag (TDRF), Timer/Event Counter 1 interrupt request flag (T1F), enable multi-function interrupt (EMFI), enable serial interface interrupt bit (ESBI), and enable Timer/Event Counter 1 interrupt bit (ET1I), on the other hand, constitute the Interrupt Control register 1 (INTC1) which is located at 1EH in the RAM.

The enable Remote control timer rising edge interrupt (ERMTO), enable Remote control timer falling edge interrupt (ERMT1), enable Remote control timer overflow interrupt (ERMTV), enable Remote control timer start counting (RME), select the Remote control timer clock source (RMCS), and select the Remote control timer clock (RMS0, RMS1) constitute the Remote timer control register (RMTC) which is located at 21 H in the RAM. EMI, EEIO, EEI1, ETOI, ET1I, ESBI, ERTI, EMFI,ERMT0 and ERMT1 are all used to control the enable/disable status of interrupts. These bits prevent the requested interrupt from being serviced. Once the interrupt request flags (MFF, TDRF, T0F, T1F, EIF1, EIF0) are all set, they remain in the INTC0~INTC1 respectively until the interrupts are serviced or cleared by a software instruction.

It is recommended that a program should not use the "CALL subroutine" within the interrupt subroutine. This is because interrupts often occur in an unpredictable manner or require to be serviced immediately in some applications. During that period, if only one stack is left, and enabling the interrupt is not well controlled, operation of the "CALL" in the interrupt subroutine may damage the original control sequence.

## Oscillator Configuration

The HT49RV9/HT49CV9 provides three oscillator circuits for system clocks, i.e., RC oscillator, crystal oscillator and 32768 Hz crystal oscillator, determined by options. No matter what type of oscillator is selected, the signal is used for the system clock. The HALT mode stops the system oscillator (RC and crystal oscillator
only) and ignores external signals so as to conserve power. The 32768 Hz crystal oscillator still runs in the HALT mode. If the 32768 Hz crystal oscillator is selected as the system oscillator, the system oscillator is not stopped but the instruction execution is stopped. Since the 32768 Hz oscillator is also designed for timing purposes, the internal timing (RTC, WDT) operation still runs even if the system enters the HALT mode.

Of the three oscillators, if the RC oscillator is used, an external resistor between OSC1 and VSS is required, and the range of the resistance should be from $24 \mathrm{k} \Omega$ to $1 \mathrm{M} \Omega$. The system clock, divided by 4 , is available on OSC2 with pull-high resistor, which can be used to synchronize external logic. The RC oscillator provides the most cost effective solution. However, the frequency of the oscillation may vary with VDD, temperature, and the chip itself due to process variations. It is therefore not suitable for timing sensitive operations where accurate oscillator frequency is desired.

On the other hand, if the crystal oscillator is selected, a crystal across OSC1 and OSC2 is needed to provide the feedback and phase shift for the oscillator, and no other external components are required. A resonator may be connected between OSC1 and OSC2 to replace the crystal and to get a frequency reference, but two external capacitors on OSC1 and OSC2 are required.
There is another oscillator circuit designed for the real time clock. In this case, only a 32768 Hz crystal oscillator can be applied. The crystal should be connected between OSC3 and OSC4.

The RTC oscillator circuit can be controlled to start-up quickly by setting the "QOSC" bit (bit 4 of RTCC). It is recommended to turn on the quick start-up function during power-on, and then turn it off after 2 seconds.

The WDT oscillator is a free running on-chip RC oscillator, and no external components are required. Although the system enters the power down mode, the system clock stops and the WDT oscillator still works with a period of approximately $65 \mu$ s at 5 V . The WDT oscillator can be disabled by options so as to conserve power.


## System Oscillator

Note: *32768Hz crystal enable condition: For WDT clock source or for system clock source.

## Watchdog Timer - WDT

The WDT clock source is implemented by a dedicated RC oscillator (WDT oscillator) or an instruction clock (system clock/4) or a real time clock oscillator (RTC oscillator). The timer is designed to prevent a software malfunction or sequence from jumping to an unknown location with unpredictable results. The WDT can be disabled by options. But if the WDT is disabled, all executions related to the WDT lead to no operation.

Once an internal oscillator (RC oscillator with a period of $65 \mu$ s at 5 V ) is selected, it is divided by $2^{12} \sim 2^{15}$ (by configuration option to get the WDT time-out period). The minimum WDT time-out period is $300 \mathrm{~ms} \sim 600 \mathrm{~ms}$. This time-out period may vary with temperature, VDD and process variations. By selecting the WDT configuration option, longer time-out periods can be realized. If the WDT time-out is selected, $2^{15}$, the maximum time-out period, divided by $2^{15} \sim 2^{16}$ is $2.1 \mathrm{~s} \sim 4.3 \mathrm{~s}$. If the WDT oscillator is disabled, the WDT clock may still come from the instruction clock and operates in the same manner except that in the halt state the WDT may stop counting and lose its protecting purposes. In this situation the logic can only be restarted by external logic. If the device operates in a noisy environment, using the on-chip RC oscillator (WDT OSC) is strongly recommended since the HALT will stop the system clock.

The WDT overflow under normal operation initializes a "chip reset" and sets the status bit "TO". In the HALT mode, the overflow initializes a "warm reset", and only the Program Counter and SP are reset to zero. To clear the contents of the WDT, there are three methods to be adopted, i.e., external reset (a low level to $\overline{\mathrm{RES}}$ ), software instruction, and a "HALT" instruction. There are two types of software instructions; "CLR WDT" and the other set - "CLR WDT1" and "CLR WDT2". Of these two types of instruction, only one type of instruction can be active at a time depending on the options - "CLR WDT" times selection option. If the "CLR WDT" is selected (i.e., CLR WDT times equal one), any execution of the "CLR WDT" instruction clears the WDT. In the case where the two "CLR WDT1" and "CLR WDT2" instructions are chosen (i.e., CLR WDT times equal two), these two instructions have to be executed to clear the WDT; otherwise, the WDT may reset the chip due to time-out.

## Multi-function Timer

The HT49RV9/HT49CV9 provides a multi-function timer for the WDT and RTC but with different time-out periods. The multi-function timer consists of an 8-stage divider and a 7-bit prescaler, with the clock source coming from the RTC OSC or the instruction clock (i.e., system clock divided by 4). The multi-function timer also provides a selectable frequency signal (ranging from $\mathrm{f}_{\mathrm{S}} / 2^{0}$ to $\mathrm{f}_{\mathrm{S}} / 2^{7}$ ) for the VFD driver circuits, and a selectable frequency signal (ranging from $\mathrm{f}_{\mathrm{S}} / 2^{1}$ to $\mathrm{f}_{\mathrm{S}} / 2^{8}$ ) for the buzzer output by options. It is recommended to select a frequency as close as possible to 32 kHz for the VFD driver circuits to obtain good display clarity.


## Real Time Clock - RTC

The real time clock (RTC) is used to supply a regular internal interrupt. Its time-out period ranges from $\mathrm{f}_{\mathrm{S}} / 2^{8}$ to $\mathrm{f}_{\mathrm{S}} / 2^{15}$ by software programming. Writing data to RT2, RT1 and RT0 (bit 2, 1, 0 of RTCC; 09H) yields various time-out periods. If the RTC time-out occurs and the interrupt is enabled, the related interrupt request flag (RTF; bit 1 of MFIS) is set.

| RT2 | RT1 | RT0 | RTC Clock Divided Factor |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | $2^{8 *}$ |
| 0 | 0 | 1 | $2^{9 *}$ |
| 0 | 1 | 0 | $2^{10 *}$ |
| 0 | 1 | 1 | $2^{11 *}$ |
| 1 | 0 | 0 | $2^{12}$ |
| 1 | 0 | 1 | $2^{13}$ |
| 1 | 1 | 0 | $2^{14}$ |
| 1 | 1 | 1 | $2^{15}$ |

Note: * not recommended to be used



The RTCC register descriptions are listed below.

| Bit No. | Label | Read/Write | Reset | Function |
| :---: | :---: | :---: | :---: | :--- |
| $0 \sim 2$ | RT0~RT2 | R/W | 1 | 8 to 1 multiplexer control inputs to select the real clock prescaler out- <br> put |
| 3 | - | - | - | Unused bit, read as "0" |
| 4 | QOSC | R/W | 0 | 32768 Hz OSC quick start-up oscillator <br> $0 / 1:$ quick/slow/start |
| 6,7 | - | - | - | Unused bit, read as "0" |

RTCC (09H) Register

## Power Down Operation - HALT

The HALT mode is initialized by the "HALT" instruction and results in the following.

- The system oscillator turns off but the WDT oscillator keeps running (if the WDT oscillator or the real time clock is selected).
- The contents of the on-chip RAM and of the registers remain unchanged.
- The WDT is cleared and starts recounting (if the WDT clock source is from the WDT oscillator or the real time clock oscillator).
- All I/O ports maintain their original status.
- The PDF flag is set but the TO flag is cleared.
- The VFD driver keeps running (if the WDT OSC or RTC OSC is selected).

The system can leave the HALT mode by means of an external reset, an interrupt, an external falling edge signal on port A, an external rising/falling edge on the RMT pin, or a WDT overflow. An external reset will initialize a chip reset and a WDT overflow will initialize a "warm reset". After examining the TO and PDF flags, the source of the reset can be determined. The PDF flag is cleared by a system power-up or by executing the "CLR WDT" instruction, and is set by executing the "HALT" instruction. The TO flag is set if a WDT time-out occurs, and causes a wake-up that only resets the Program Counter and SP, the other flags remain in their original status.

The port A wake-up and interrupt methods can be considered as a continuation of normal execution. Each bit in port A can be independently selected to wake up the device by options. Awakening from an I/O port stimulus, the program will resume execution at the next instruction. If the system is woken up via an interrupt, two possibilities may occur. If the related interrupt is disabled or the interrupt is enabled but the stack is full, the
program will resume execution at the next instruction. But if the interrupt is enabled and the stack is not full, the regular interrupt response takes place.

If an interrupt request flag is set to " 1 " before entering the "HALT" mode, the wake-up function of the related interrupt will be disabled.

If a wake-up event occurs, it takes $1024 \mathrm{t}_{\mathrm{SYS}}$ (system clock period) to resume normal operation. In other words, a dummy period will be inserted after a wake-up. If the wake-up results from an interrupt acknowledge signal, the actual interrupt subroutine execution will be delayed by one or more cycles. However, if the wake-up results in the next instruction execution, this will be executed immediately after the dummy period is finished.

To minimize power consumption, all the I/O pins should be carefully managed before entering the HALT mode.

## Reset

There are three ways in which a microcontroller reset can occur, through events occurring both internally and externally:

- $\overline{\mathrm{RES}}$ is reset during normal operation
- $\overline{R E S}$ is reset during HALT
- WDT time-out is reset during normal operation

The WDT time-out reset during HALT is a little different from other kinds of reset. Most of the conditions remain unchanged except that the program counter and stack pointer will be cleared to 0 and the TO flag will be set to 1. Most registers are reset to the "initial condition" once the reset conditions are met.

The different types of resets described affect the reset flags in different ways. These flags, the PDF and TO flags, are located in the status register and are con-
trolled by various microcontroller operations such as the HALT function or Watchdog Timer. The reset flags are shown in the table:

| TO | PDF | RESET Conditions |
| :---: | :---: | :--- |
| 0 | 0 | $\overline{R E S}$ reset during power-up |
| $u$ | $u$ | $\overline{R E S}$ reset during normal operation |
| 0 | 1 | $\overline{R E S}$ Wake-up HALT |
| 1 | $u$ | WDT time-out during normal operation |
| 1 | 1 | WDT Wake-up HALT |

Note: "u" stands for unchanged
To ensure that the system oscillator is started and stabilized, the SST (System Start-up Timer) provides an extra delay of 1024 system clock pulses when the system awakes from the HALT state or during power-on.
An extra SST delay is added during the power-on period, and any wake-up from HALT may enable only the SST delay.

The functional unit chip reset status is shown below.

| Program Counter | 000 H |
| :--- | :--- |
| Interrupt | Disabled |
| Prescaler, Divider | Cleared |
| Remote control timer | Cleared. After a master reset, <br> WDT starts counting |
| Timer/event Counter | Off |
| Input/output Ports | Input mode |
| Stack Pointer | Points to the top of the stack |



## Reset Circuit

Note: "**" Make the length of the wiring, which is connected to the $\overline{R E S}$ pin as short as possible, to avoid noise interference.


Reset Timing Chart


Reset Configuration

The register states are summarized below:

| Register | Reset (Power-on) | WDT Time-out (Normal Operation) | $\overline{\text { RES }}$ Reset (Normal Operation) | $\overline{\text { RES }}$ Reset (HALT) | WDT Time-out (HALT)* |
| :---: | :---: | :---: | :---: | :---: | :---: |
| TMROH | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuuu uuuu |
| TMROL | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | unuu uuuu |
| TMR0C | 00-0 1000 | 00-0 1000 | 00-0 1000 | 00-0 1000 | uu-u uuuu |
| TMR1H | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuuu uuuu |
| TMR1L | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuuu uuuu |
| TMR1C | 0000 1--- | 0000 1--- | 0000 1--- | 0000 1--- | uuuu u--- |
| Program Counter | 0000H | 0000H | 0000H | 0000H | 0000H |
| MP0 | xxxx xxxx | uuuu uuuu | uuuu uuuu | uuuu uuuu | uuun unuu |
| MP1 | xxxx xxxx | uuuu uuuu | uuuu uuuu | uuuu uuuu | uuuu unuu |
| BP | 00000000 | 00000000 | 00000000 | 00000000 | uuuu uuuu |
| ACC | xxxx xxxx | unuu unuu | uuuu uuuu | uuuu uuuu | uuuu uuuu |
| TBLP | xxxx xxxx | uuuu unuu | uuuu uuuu | uuuu uuuu | uuuu unuu |
| TBLH | xxxx xxxx | uuuu unuu | uuuu uuuu | uuuu uuuu | uuuu uuuu |
| STATUS | --00 xxxx | --1u uuuu | --uu uuuu | --01 uuuu | --11 uuuu |
| INTC0 | 00000000 | 00000000 | 00000000 | 00000000 | uuuu uuuu |
| INTC1 | -000-000 | -000-000 | -000-000 | -000-000 | -uuu -uuu |
| RMTC | 0000 000- | 0000 000- | 0000 000- | 0000 000- | uuuu uuu- |
| MFIS | ---0 0000 | ---0 0000 | ---0 0000 | ---0 0000 | ---u uuuu |
| RTCC | ---0-111 | ---0-111 | ---0-111 | ---0-111 | ---u -uuu |
| PA | 11111111 | 11111111 | 11111111 | 11111111 | uuun uuuu |
| PAC | 11111111 | 11111111 | 11111111 | 11111111 | uuuu unuu |
| PB | 11111111 | 11111111 | 11111111 | 11111111 | unuu unuu |
| PBC | 11111111 | 11111111 | 11111111 | 11111111 | uuun unuu |
| PC | 11111111 | 11111111 | 11111111 | 11111111 | uuuu uuuu |
| PCC | 11111111 | 11111111 | 11111111 | 11111111 | uuuu uuuu |
| PD | 11111111 | 11111111 | 11111111 | 11111111 | uuun unuu |
| PDC | 11111111 | 11111111 | 11111111 | 11111111 | unuu unuu |
| PWMO | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | unuu unuu |
| PWM1 | xxxx xxxx | xxxx xxxx | xxxx xxxx | $x x x x$ xxxx | uuuu uuuu |
| PWM2 | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuun unuu |
| PWM3 | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuuu uuuu |
| SBCR | 01100000 | 01100000 | 01100000 | 01100000 | unuu unuu |
| SBDR | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuun unuu |
| RMT0 | 00000000 | 00000000 | 00000000 | 00000000 | uuuu unuu |
| RMT1 | 00000000 | 00000000 | 00000000 | 00000000 | uuuu uuuu |
| ADRL | xx-- ---- | xx------ | xx-- ---- | xx-- ---- | uu-- ---- |
| ADRH | xxxx xxxx | xxxx xxxx | xxxx xxxx | xxxx xxxx | uuun unuu |
| ADCR | 01000000 | 01000000 | 01000000 | 01000000 | uuun uuuu |
| ACSR | 1----00 | 1----00 | 1----00 | ------00 | u--- --uu |
| VFDC | 00001111 | 00001111 | 00001111 | 00001111 | uuuu uuuu |

Note: "*" stands for warm reset
"u" stands for unchanged
" x " stands for unknown

## Timer/Event Counter

Two timer/event counters (TMR0,TMR1) are implemented in the microcontroller. The Timer/Event Counter 0 contains a 16-bit programmable count-up counter and the clock may come from an external source or an internal clock source. An internal clock source comes from $\mathrm{f}_{\mathrm{SY}}$. The Timer/Event Counter 1 contains a 16-bit programmable count-up counter and the clock may come from an external source or an internal clock source. An internal clock source comes from $\mathrm{f}_{\mathrm{SYS}} / 4$ or 32768 Hz selected by option. The external clock input allows the user to count external events, measure time intervals or pulse widths.

There are six registers related to the Timer/Event Counter 0; TMROH (OCH), TMROL (ODH), TMROC (OEH) and the Timer/Event Counter 1, TMR1H (OFH), TMR1L (10H), and TMR1C (11H). Writing TMR0L (TMR1L) will only place the written data to an internal lower-order byte buffer (8-bit) and writing TMROH (TMR1H) will transfer the specified data and the contents of the lower-order byte buffer to TMR0H (TMR1H) and TMR0L (TMR1L) registers, respectively. The Timer/Event Counter 1/0 preload register is changed by each writing TMR0H (TMR1H) operations. Reading TMROH
(TMR1H) will latch the contents of TMR0H (TMR1H) and TMROL (TMR1L) counters to the destination and the lower-order byte buffer, respectively. Reading the TMR0L (TMR1L) will read the contents of the lower-order byte buffer. The TMR0C (TMR1C) is the Timer/Event Counter 0 (1) control register, which defines the operating mode, counting enable or disable and an active edge.

The TM0 and TM1 bits define the operation mode. The event count mode is used to count external events, which means that the clock source is from an external (TMR0, TMR1) pin. The timer mode functions as a normal timer with the clock source coming from the internal selected clock source. Finally, the pulse width measurement mode can be used to count the high or low level duration of the external signal (TMR0, TMR1), and the counting is based on the internal selected clock source.

In the event count or timer mode, the timer/event counter starts counting at the current contents in the timer/ event counter and ends at FFFFH. Once an overflow occurs, the counter is reloaded from the timer/event counter preload register, and generates an interrupt request flag (TOF; bit 6 of INTC0, T1F; bit 4 of INTC1). In the pulse width measurement mode with the values of the


Timer/Event Counter 0


Timer/Event Counter 1

TON and TE bits equal to 1, after the TMR0 (TMR1) has received a transient from low to high (or high to low if the TE bit is " 0 "), it will start counting until the TMR0 (TMR1) returns to the original level and resets the TON. The measured result remains in the timer/event counter even if the activated transient occurs again. In other words, only 1-cycle measurement can be made until the TON bit is set. The cycle measurement will continue as long as it receives further transient pulse. In this operation mode, the timer/event counter begins counting not according to the logic level but to the transient edges. In the case of counter overflows, the counter is reloaded
from the timer/event counter register and issues an interrupt request, as in the other two modes, i.e., the event and timer modes.

To enable the counting operation, the Timer on bit (TON; bit 4 of TMR0C or TMR1C) should be set to 1 . In the pulse width measurement mode, TON is automatically cleared after the measurement cycle is completed. But in the other two modes, TON can only be reset by instructions. The overflow of the Timer/Event Counter 0/1 is one of the wake-up sources and can also be applied to a PFD (Programmable Frequency Divider) output at

| Bit No. | Label | Function |
| :---: | :---: | :---: |
| 0~2 | TOPSC0~ TOPSC2 | Defines the prescaler stages TOPSC2, TOPSC1, TOPSC0= $000: f_{I N T}=f_{S Y S}$ <br> 001: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 2$ <br> 010: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 4$ <br> 011: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 8$ <br> 100: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 16$ <br> 101: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 32$ <br> 110: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 64$ <br> 111: $\mathrm{f}_{\mathrm{INT}}=\mathrm{f}_{\mathrm{SYS}} / 128$ |
| 3 | T0E | Defines the TMRO active edge of the timer/event counter: <br> In Event Counter Mode (TOM1,TOM0) $=(0,1)$ : <br> 1:count on falling edge; <br> 0:count on rising edge <br> In Pulse Width measurement mode (TOM1,TOM0)=(1,1): <br> 1: start counting on the rising edge, stop on the falling edge; <br> 0 : start counting on the falling edge, stop on the rising edge |
| 4 | T0ON | Enable/disable timer counting ( $0=$ disable; $1=$ enable) |
| 5 | - | Unused bit, read as "0" |
| $\begin{aligned} & 6 \\ & 7 \end{aligned}$ | $\begin{aligned} & \text { T0M0 } \\ & \text { T0M1 } \end{aligned}$ | $\begin{aligned} & \text { Defines the operating mode (TOM1, TOM0) } \\ & 01=\text { Event count mode (External clock) } \\ & 10=\text { Timer mode (Internal clock) } \\ & 11=\text { Pulse Width measurement mode (External clock) } \\ & 00=\text { Unused } \end{aligned}$ |

TMROC (0EH) Register

| Bit No. | Label | Function |
| :---: | :---: | :---: |
| 0~2 | - | Unused bit, read as "0" |
| 3 | T1E | Defines the TMR1 active edge of the timer/event counter: <br> In Event Counter Mode (T1M1,T1M0)=(0,1): <br> 1:count on falling edge; <br> 0:count on rising edge <br> In Pulse Width measurement mode (T1M1,T1M0)=(1,1): <br> 1: start counting on the rising edge, stop on the falling edge; <br> 0 : start counting on the falling edge, stop on the rising edge |
| 4 | T1ON | Enable/disable timer counting ( $0=$ disable; $1=$ enable) |
| 5 | T1S | Defines the TMR1 internal clock source ( $0=\mathrm{f}_{\text {SYS }} / 4 ; 1=32768 \mathrm{~Hz}$ ) |
| $\begin{aligned} & 6 \\ & 7 \end{aligned}$ | $\begin{aligned} & \text { T1M0 } \\ & \text { T1M1 } \end{aligned}$ | Defines the operating mode (T1M1, T1M0) <br> 01= Event count mode (External clock) <br> 10= Timer mode (Internal clock) <br> $11=$ Pulse Width measurement mode (External clock) <br> 00= Unused |

TMR1C (11H) Register

PA3 by options. Only one PFD (PFD0 or PFD1) can be applied to PA3 by options. No matter what the operation mode is, writing a 0 to ETOI or ET1I disables the related interrupt service. When the PFD function is selected, executing "SET [PA].3" instruction will enable the PFD output and executing "CLR [PA].3" instruction will disable the PFD output.

In the case of timer/event counter OFF condition, writing data to the timer/event counter preload register also reloads that data to the timer/event counter. But if the timer/event counter is turned on, data written to the timer/event counter is kept only in the timer/event counter preload register. The timer/event counter still continues its operation until an overflow occurs.

When the timer/event counter (reading TMR0/TMR1) is read, the clock is blocked to avoid errors, as this may results in a counting error. Blocking of the clock should be taken into account by the programmer. It is strongly recommended to load a desired value into the TMR0/TMR1 register first, before turning on the related timer/event
counter, for proper operation since the initial value of TMR0/TMR1 is unknown. Due to the timer/event counter scheme, the programmer should pay special attention on the instruction to enable then disable the timer for the first time, whenever there is a need to use the timer/event counter function, to avoid unpredictable result. After this procedure, the timer/event function can be operated normally.

Bits 2~0 of the TMROC can be used to define the pre-scaling stages of the internal clock sources of the timer/event counter. The definitions are as shown. The overflow signal of the timer/event counter can be used to generate the PFD signal. The timer prescaler is also used as the PWM counter.

## Remote Control Timer - RMT

The HT49RV9/HT49CV9 provides an 8-bit remote control timer that has a pulse width measurement function. Pulse width is measured from a difference in count value when the valid edge (RMT pin) has been detected while the timer operates in the running mode.


Remote Control Timer

| Bit No. | Label | Function |
| :---: | :---: | :--- |
| 0 | - | Unused bit, read as " $0^{\prime \prime}$ |$|$| 1 | ERMT0 | Controls the remote control timer rising edge interrupt (1=enable; $0=$ disable) |
| :---: | :---: | :--- |
| 2 | ERMT1 | Controls the remote control timer falling edge interrupt (1=enable; $0=$ disable) |

RMTC (21H) Register
The RMT pin with rising/falling edge wake-up function, 8-bit timer counter will be cleared during a chip reset.

## Serial Interface

The Serial Interface function has four basic signals included. They are the SDI (serial data input), SDO (serial data output), SCK (serial clock) and $\overline{\text { SCS }}$ (slave select pin).

Two registers (SBCR \& SBDR) unique to the serial interface provide control, status and data storage.


SIO Timing (SIOCLK Configuration is Rising Edge)

| Bit No. | Label |  |
| :---: | :---: | :--- |
| 0 | - | Unused bit, read as " 0 " |$\quad$ Function

## SBCR (1FH) Register

- SBCR: Serial bus control register
- Bit 7 (CKS): clock source selection:
$\mathrm{f}_{\mathrm{SIO}}=\mathrm{f}_{\mathrm{SYS}} / 4$ or $\mathrm{f}_{\text {RTCOSC }}$
- Bit 6 (M1), Bit 5 (M0): master/slave mode \& baud rate selection
- M1, M0

00: master mode, baud rate $=f_{\text {SIO }}$
01: master mode, baud rate $=f_{\text {sIo }} / 4$
10: master mode, baud rate $=\mathrm{f}_{\text {SIO }} / 16$
11: slave mode

- Bit 4 (SBEN): serial bus enable/disable (1/0)
- Enable: ( $\overline{\text { SCS }}$ dependent on CSEN bit) Disable $\rightarrow$ enable: SCK, SDI, SDO, $\overline{\text { SCS }}=0$ and waiting for writing data to SBDR (TXRX buffer) Master mode: write data to SBDR (TXRX buffer) $\rightarrow$ start transmission/reception automatically Master mode: when data has been transferred $\rightarrow$ set TRF
Slave mode: when a SCK (and $\overline{\text { SCS }}$ dependent on CSEN) is received, data in TXRX buffer is shifted-out and data on SDI is shifted-in
- Disable: SCK, SDI, $\overline{\text { SCS }}$ floating, SDO output high
- Bit 3 (MLS): MSB or LSB (1/0) shift first control bit
- Bit 2 (CSEN): serial bus selection signal enable/disable ( $\overline{\mathrm{SCS}}$ ), when CSEN $=0, \overline{\mathrm{SCS}}$ is floating
- Bit 1 (WCOL): this bit is set to 1 if data is written to SBDR (TXRX buffer) when data is transferred $\rightarrow$ writing will be ignored if data is written to SBDR (TXRX buffer) when data is transferred
- Bit 0 (TRF): data transferred or data received $\rightarrow$ used to generate interrupt
Note: data receiving is still working when the MCU enters the HALT mode
- SBDR: Serial bus data register
- Data written to SBDR $\rightarrow$ write data to TXRX buffer only
- Data read from SBDR $\rightarrow$ read from SBDR only
- Operating Mode description
- Master transmitter: clock sending and data I/O started by writing SBDR
- Master clock sending started by writing SBDR
- Slave transmitter: data I/O started by clock received
- Slave receiver: data I/O started by clock received


## Clock Polarity=Rising Edge or Falling Edge (Configuration Option)

- Serial interface operation
- Master mode operation

Step1: Select CKS and select M1, M0 $=00,01,10$
Step2: Select CSEN, MLS (the same as the slave)
Step3: Set SBEN
Step4: Writing data to SBDR

- data is stored in TXRX buffer
- output SCK and $\overline{\text { SCS }}$ signals
- go to step 5

Note: SIO internal operation:

* data stored in TXRX buffer, and SDI data is shifted into TXRX buffer
* data transferred, data in TXRX buffer is latched into SBDR
Step5: Check WCOL
- WCOL=1, clear WCOL and go to step 4
- WCOL= 0, go to step 6

Step6: Check TRFor waiting for serial bus interrupt Step7: Read data from SBDR
Step8: Clear TRF
Step9: Go to step 4


SIO Block Diagram

- Slave mode operation

Step1: CKS don't care and select M1, M0 $=11$
Step2: Select CSEN, MLS (the same as the master) Step3: Set SBEN
Step4: Writing data to SBDR

- data is stored in the TXRX buffer
- waiting for master clock signal (and $\overline{\text { SCS }}$ ): SCK
- go to step 5

Note: SIO internal operation:

* SCK (SCS) received
* output data in TXRX buffer and SDI data is shifted into TXRX buffer
* data transferred, data in TXRX buffer is latched into SBDR
Step5: Check WCOL
- WCOL=1, clear WCOL, go to step 4
- WCOL=0, go to step 6

Step6: Check TRFor waiting for serial bus interrupt
Step7: Read data from SBDR
Step8: Clear TRF
Step9: Go to step 4

## Input/Output Ports

There are 32 bidirectional input/output lines in the microcontroller, labeled as PA, PB, PC and PD, which are mapped to the data memory of [12H], [14H], [16H] and $[18 \mathrm{H}]$, respectively. All of these I/O ports can be used for input and output operations. For input operation, these ports are non-latching, that is, the inputs must be ready at the T2 rising edge of instruction "MOV $\mathrm{A},[\mathrm{m}]^{\prime \prime}(\mathrm{m}=12 \mathrm{H}, 14 \mathrm{H}, 16 \mathrm{H}$ or 18 H ). For output operation, all the data is latched and remains unchanged until the output latch is rewritten.
Each I/O line has its own control register (PAC, PBC, PCC, PDC) to control the input/output configuration. With this control register, CMOS output or Schmitt Trigger input with or without pull-high resistor structures can be reconfigured dynamically under software control. To function as an input, the corresponding latch of the control register must write a " 1 ". The input source also depends on the control register. If the control register bit is " 1 ", the input will read the pad state. If the control register bit is " 0 ", the contents of the latches will move to the internal bus. The latter is possible in the "read-modifywrite" instruction.

For an output function, CMOS is the only configuration. These control registers are mapped to locations 13 H , $15 \mathrm{H}, 17 \mathrm{H}$ and 19 H .

After a chip reset, these input/output lines remain at high levels or floating state (depending on pull-high options). Each bit of these input/output latches can be set or cleared by "SET [m].i" and "CLR [m].i" (m=12H, 14H, 16 H or 18 H ) instructions.
Some instructions first input data and then follow the output operations. For example, "SET [m].i", "CLR
[m].i", "CPL [m]", "CPLA [m]" read the entire port states into the CPU, execute the defined operations (bit-operation), and then write the results back to the latches or the accumulator.

Each line of port A has the capability of waking-up the device. Each I/O port has a pull-high option. Once the pull-high option is selected, the I/O port has a pull-high resistor, otherwise, there's none. Take note that a non-pull-high I/O port operating in input mode will cause a floating state.

The PA3 pin is pin-shared with the PFD signal. If the PFD option is selected, the output signal in the output mode of PA3 will be the PFD signal generated by the timer/event counter overflow signal. The input mode always retains its original functions. Once the PFD option is selected, the PFD output signal is controlled by the PA3 data register only. Writing a "1" to the PA3 data register will enable the PFD output function and writing a " 0 " will force the PA3 pin to remain at " 0 ". The I/O functions of PA3 are shown below.

| I/O <br> Mode | I/P <br> (Normal) | O/P <br> (Normal) | I/P <br> (PFD) | O/P <br> (PFD) |
| :---: | :---: | :---: | :---: | :---: |
| PA3 | Logical <br> Input | Logical <br> Output | Logical <br> Input | PFD <br> (Timer on) $)$ |

Note: The PFD frequency is the timer/event counter overflow frequency divided by 2 .

The PA0, PA1, PA3, PD4, PD5, PD6 and PD7 pins are pin-shared with the BZ, $\overline{\mathrm{BZ}}, \mathrm{PFD}, \overline{\mathrm{INTO}}, \overline{\mathrm{INT} 1}, \mathrm{TMRO}$ and TMR1 pins, respectively.
The descriptions of PFD control signal and PFD output frequency are listed in the following table.

| Timer | Timer <br> Preload <br> Value | PA3 Data <br> Register | PA3 Pad <br> State | PFD <br> Frequency |
| :---: | :---: | :---: | :---: | :---: |
| OFF | X | 0 | 0 | X |
| OFF | X | 1 | U | X |
| ON | N | 0 | 0 | X |
| ON | N | 1 | PFD | $\mathrm{f}_{\mathrm{TMR} /[2 \times(\mathrm{M}-\mathrm{N})]}$ |

Note: "X" stands for unused
"U" stands for unknown
" M " is " 65536 " for PFD0 or PFD1
" N " is preload value for timer/event counter
"fTMR" is input clock frequency for timer/event counter

The PA0 and PA1 pins are pin-shared with the BZ and $\overline{B Z}$ signal, respectively. If the $B Z / \overline{B Z}$ option is selected, the output signal in the output mode of PAO/PA1 will be the buzzer signal generated by the multi-function timer. The input mode always remains in its original function. Once the $B Z / \overline{B Z}$ option is selected, the buzzer output signal are controlled by the PA0/PA1 data register only.


## Input/Output Ports

The I/O function of PA0/PA1 are shown below.

| PA0 I/O | I | I | O | O | O | O | O | O | O | O |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| PA1 I/O | I | O | I | I | I | O | O | O | O | O |
| PA0 Mode | X | X | C | B | B | C | B | B | B | B |
| PA1 Mode | X | C | X | X | X | C | C | C | B | B |
| PA0 Data | X | X | D | 0 | 1 | $D_{0}$ | 0 | 1 | 0 | 1 |
| PA1 Data | X | D | X | X | X | D1 | D | D | X | X |
| PA0 Pad Status | I | I | D | 0 | B | $D_{0}$ | 0 | B | 0 | B |
| PA1 Pad Status | I | D | I | I | I | $D_{1}$ | D | D | 0 | B |

Note: "I" input; "O" output
"D, D0, D1" Data
"B" buzzer option, BZ or $\overline{B Z}$
" X " don't care
"C" CMOS output
The PB port can also be used as A/D converter inputs. The A/D function will be described later. There is a PWM function shared with PD0/PD1/PD2/PD3. If the PWM function is enabled, the PWM0/PWM1/PWM2/PWM3 signal will appear on PD0/PD1/PD2/PD3 (if PD0/PD1/ PD2/PD3 is operating in the output mode). Writing a "1" to PD0~PD3 data register will enable the PWM output function and writing a " 0 " will force the PD0~PD3 to remain at " 0 ". The I/O functions of PD0/PD1/PD2/PD3 are as shown.

| I/O <br> Mode | I/P <br> (Normal) | O/P <br> (Normal) | I/P <br> (PWM) | O/P <br> (PWM) |
| :---: | :---: | :---: | :---: | :---: |
| PD0 |  |  |  | PWM0 |
| PD1 | Logical | Logical | Logical | PWM1 |
| PD2 | Input | Output | Input | PWM2 |
| PD3 |  |  |  | PWM3 |

It is recommended that unused or not bonded out I/O lines should be set as output pins by software instruction to avoid consuming power under input floating state.

## PWM

The microcontroller provides a 4 channel and $(6+2) /(7+1)$ (dependent on options) bits PWM output shared with PD0/PD1/PD2/PD3. The PWM channels have their data registers denoted as PWM0 (1AH), PWM1 (1BH), PWM2 (1CH) and PWM3 (1DH). The frequency source of the PWM counter comes from fsYs. The PWM registers are two 8-bit registers. The waveforms of the PWM outputs are as shown. Once the PD0/PD1/PD2/PD3 are selected as the PWM outputs and the output function of PD0/PD1/PD2/PD3 are enabled (PDC.0/PDC.1/ PDC.2/PDC.3="0"), writing a "1" to PD0/PD1/PD2/PD3 data register will enable the PWM output function and writing a " 0 " will force the PD0/PD1/PD2/PD3 to remain at " 0 ".
A (6+2) bits mode PWM cycle is divided into four modulation cycles (modulation cycle 0~modulation cycle 3). Each modulation cycle has 64 PWM input clock period. In a (6+2) bit PWM function, the contents of the PWM register is divided into two groups. Group 1 of the PWM register is denoted by DC which is the value of PWM.7~PWM.2.

The group 2 is denoted by AC which is the value of PWM.1~PWM.0.

In a (6+2) bits mode PWM cycle, the duty cycle of each modulation cycle is shown in the table.

| Parameter | AC (0~3) | Duty Cycle |
| :---: | :---: | :---: |
| Modulation cycle i <br> $(\mathrm{i}=0 \sim 3)$ | $\mathrm{i}<\mathrm{AC}$ | $\frac{\mathrm{DC}+1}{64}$ |
|  | $\mathrm{i} \geq \mathrm{AC}$ | $\frac{\mathrm{DC}}{64}$ |

A (7+1) bits mode PWM cycle is divided into two modulation cycles (modulation cycle0~modulation cycle1). Each modulation cycle has 128 PWM input clock period.

In a (7+1) bits PWM function, the contents of the PWM register is divided into two groups. Group 1 of the PWM register is denoted by DC which is the value of PWM.7~PWM. 1.

The group 2 is denoted by AC which is the value of PWM.0.

In a (7+1) bits mode PWM cycle, the duty cycle of each modulation cycle is shown in the table.

| Parameter | AC (0~1) | Duty Cycle |
| :---: | :---: | :---: |
| Modulation cycle i <br> $(\mathrm{i}=0 \sim 1)$ | $\mathrm{i}<\mathrm{AC}$ | $\frac{\mathrm{DC}+1}{128}$ |
|  | $\mathrm{i} \geq \mathrm{AC}$ | $\frac{\mathrm{DC}}{128}$ |

The modulation frequency, cycle frequency and cycle duty of the PWM output signal are summarized in the following table.

| PWM <br> Modulation Frequency | PWM Cycle <br> Frequency | PWM Cycle <br> Duty |
| :---: | :---: | :---: |
| $\mathrm{f}_{\text {SYS }} / 64$ for $(6+2)$ bits mode <br> $\mathrm{f}_{\text {SYS }} / 128$ for (7+1) bits mode | $\mathrm{f}_{\mathrm{SYS}} / 256$ | $[\mathrm{PWM}] / 256$ |


(6+2) PWM Mode

(7+1) PWM Mode

## A/D Converter

The 8 channels and 10 bit resolution A/D (9 bit accuracy) converter are implemented in these microcontrollers. The reference voltage is VDD. The A/D converter contains 4 special registers which are; ADRL $(24 \mathrm{H})$, ADRH $(25 \mathrm{H}), \operatorname{ADCR}(26 \mathrm{H})$ and ACSR $(27 \mathrm{H})$. The ADRH and ADRL contain the A/D result register higher-order byte and lower-order byte and are read-only. After the A/D conversion is completed, the ADRH and ADRL should be read to retrieve the conversion result data. The ADCR is an A/D converter control register, which defines the A/D channel number, analog channel select, start A/D conversion control bit and the end of $A / D$ conversion flag. If the user wishes to start an A/D conversion, they should define the PB configuration, select the converted analog channel, and give the START bit a rising edge and falling edge $(0 \rightarrow 1 \rightarrow 0)$. At the end of the A/D conversion, the EOCB bit is cleared. The ACSR is the A/D clock setting register, which is used to select the A/D clock source.

The A/D converter control register is used to control the A/D converter. The bit2~bit0 of the ADCR are used to select an analog input channel. There are a total of eight channels to select. Bits 5~3 of the ADCR are used to set the PB configurations. PB can be an analog input or digital I/O line determined by these 3 bits. Once a PB line is selected as an analog input, the I/O functions and pull-high resistor of this I/O line are disabled and the A/D converter circuit is powered-on. The EOCB bit (bit 6 of the ADCR) is the end of $A / D$ conversion flag. Check this bit to know when the A/D conversion is completed. The START bit of the ADCR is used to begin the conversion of the A/D converter. Giving the START bit a rising edge and falling edge means that the A/D conversion has started. In order to ensure that the A/D conversion is completed, the START should remain at " 0 " until the EOCB is cleared to " 0 " (end of A/D conversion).

Bit 7 of the ACSR is used for testing purposes only. It cannot be used by the user. The bit 1 and bit 0 of the ACSR are used to select the A/D clock sources.

| Bit No. | Label |  |
| :---: | :---: | :--- |
|  |  | Selects the A/D converter clock source <br> $00=$ system clock/2 <br> 0 |
| ADCS0 | Function |  |
| 1 | ADCS1 | $10=$ system clock/8 <br> $11=$ undem clock/32 |
| $2 \sim 6$ | - | Unused bit, read as "0" |
| 7 | TEST | For test mode used only |

ACSR (27H) Register

| Bit No. | Label | Function |
| :---: | :---: | :--- |
| 0 | ACS0 |  |
| 1 | ACS1 | Defines the analog channel select. |
| 2 | ACS2 |  |
| 3 | PCR0 | Defines the port B configuration select. If PCRO, PCR1 and PCR2 are all zero, the ADC cir- |
| 4 | PCR1 | cuit is powered- off to reduce power consumption |
| 5 | PCR2 |  |
| 6 | EOCB | Provides a response at the end of the A/D conversion. ( $0=$ end of A/D conversion $)$ |
| 7 | START | Starts the A/D conversion. $(0 \rightarrow 1 \rightarrow 0=$ start; $0 \rightarrow 1=$ reset A/D converter $)$ |

ADCR (26H) Register

| ACS2 | ACS1 | ACS0 | Analog Channel |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | A0 |
| 0 | 0 | 1 | A1 |
| 0 | 1 | 0 | A2 |
| 0 | 1 | 1 | A3 |
| 1 | 0 | 0 | A4 |
| 1 | 0 | 1 | A5 |
| 1 | 1 | 0 | A6 |
| 1 | 1 | 1 | A7 |

Analog Input Channel Selection

The EOCB bit is set to " 1 " when the START bit is set from " 0 " to " 1 ".

| Register | Bit7 | Bit6 | Bit5 | Bit4 | Bit3 | Bit2 | Bit1 | Bit0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| ADRL | D1 | D0 | - | - | - | - | - | - |
| ADRH | D9 | D8 | D7 | D6 | D5 | D4 | D3 | D2 |

Note: $\quad$ D0~D9 is the A/D conversion result data bit LSB~MSB.

| PCR2 | PCR1 | PCR0 | $\mathbf{7}$ | $\mathbf{6}$ | $\mathbf{5}$ | $\mathbf{4}$ | $\mathbf{3}$ | $\mathbf{2}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | PB7 | PB6 | PB5 | PB4 | PB3 | PB2 | PB1 | PB0 |
| 0 | 0 | 1 | PB7 | PB6 | PB5 | PB4 | PB3 | PB2 | PB1 | AN0 |
| 0 | 1 | 0 | PB7 | PB6 | PB5 | PB4 | PB3 | PB2 | AN1 | AN0 |
| 0 | 1 | 1 | PB7 | PB6 | PB5 | PB4 | PB3 | AN2 | AN1 | AN0 |
| 1 | 0 | 0 | PB7 | PB6 | PB5 | PB4 | AN3 | AN2 | AN1 | AN0 |
| 1 | 0 | 1 | PB7 | PB6 | PB5 | AN4 | AN3 | AN2 | AN1 | AN0 |
| 1 | 1 | 0 | PB7 | PB6 | AN5 | AN4 | AN3 | AN2 | AN1 | AN0 |
| 1 | 1 | 1 | AN7 | AN6 | AN5 | AN4 | AN3 | AN2 | AN1 | AN0 |

Port B Configuration

The following two programming examples illustrate how to setup and implement an A/D conversion. In the first example, the method of polling the EOCB bit in the ADCR register is used to detect when the conversion cycle is completed, whereas in the second example, the $A / D$ interrupt is used to determine when the conversion is completed.
Example: using EOCB Polling Method to detect end of conversion

| clr | INTC0.7 | ; disable A/D interrupt in interrupt control register |
| :--- | :--- | :--- |
| mov | a,00100000B |  |
| mov | ADCR,a | ; setup ADCR register to configure Port PB0~PB3 as A/D inputs and select |
| mov | a,00000001B | ; AN0 to be connected to the A/D converter |
| mov | ACSR,a | ; setup the ACSR register to select $f_{S Y S} / 8$ as the A/D clock |

Start conversion:

| clr | ADCR. 7 |  |
| :--- | :--- | :--- |
| set | ADCR. 7 | ; reset A/D |
| clr | ADCR. 7 | ; start A/D |

## Polling_EOC:

sz ADCR. 6
jmp polling_EOC
mov a,ADRH
mov adrh_buffer
(
; continue polling
; read conversion result from the high byte ADRH register
, save result to user defined register
mov adrl_buffer,a ; save result to user defined register
jmp start_conversion
; start next A/D conversion


Note: A/D clock must be fsys/2, fsys/8 or fsys/32
A/D Conversion Timing

## VFD Display Memory

The HT49RV9/HT49CV9 provides an area of embedded data memory for the VFD display. This area is located from 40 H to 6FH of the RAM at Bank 1. The Bank Pointer (BP; located at 04H of the RAM) is the switch for the RAM and the VFD display memory. When the BP is set as " 1 ", any data written into $40 \mathrm{H} \sim 6 \mathrm{FH}$ will affect the VFD display. When the BP is written as " 0 ", " 2 ", " 3 " or " 4 " any data written into $40 \mathrm{H} \sim 6 \mathrm{FH}$ is meant to access
the general purpose data memory. The VFD display memory can be read and written to only by an indirect addressing mode using MP1. When data is written into the display data area, it is automatically read by the VFD driver which then generates the corresponding VFD driving signals. To turn the display on or off, a "1" or a " 0 " is written to the corresponding bit of the display memory, respectively. The figure illustrates the mapping between the display memory and VFD pattern for the HT49RV9/HT49CV9.

|  | SEG19~SEG16 | SEG15~SEG12 | SEG11~SEG8 | $\begin{gathered} \text { SEG7~ } \\ \text { SEG4 } \end{gathered}$ | $\begin{gathered} \text { SEG3~ } \\ \text { SEG0 } \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| GRID0 | 42 H | $41 \mathrm{H}_{\mathrm{U}}$ | $41 \mathrm{H}_{\mathrm{L}}$ | $40 \mathrm{H}_{\mathrm{U}}$ | $40 \mathrm{H}_{\text {L }}$ |
| GRID1 | 45 H L | $44 \mathrm{H}_{\mathrm{U}}$ | $44 \mathrm{H}_{\mathrm{L}}$ | $43 \mathrm{H}_{\mathrm{U}}$ | $43 \mathrm{H}_{\text {L }}$ |
| GRID2 | $48 \mathrm{H}_{\mathrm{L}}$ | 47 HU | 47 H | $46 \mathrm{H}_{\mathrm{u}}$ | $46 \mathrm{H}_{\mathrm{L}}$ |
| GRID3 | 4 BH L | 4 AHU | 4 AH L | $49 \mathrm{H}_{\mathrm{U}}$ | $49 \mathrm{H}_{\text {L }}$ |
| GRID4 | 4 EH L | 4 DH | 4DHL | $4 \mathrm{CH}_{U}$ | $4 \mathrm{CH}_{\mathrm{L}}$ |
| GRID5 | $51 \mathrm{H}_{\mathrm{L}}$ | $50 \mathrm{H}_{U}$ | $50 \mathrm{H}_{\mathrm{L}}$ | 4 FH | $4 \mathrm{FH}_{\text {L }}$ |
| GRID6 | $54 \mathrm{H}_{\mathrm{L}}$ | $53 \mathrm{H}_{\mathrm{U}}$ | $53 \mathrm{H}_{\mathrm{L}}$ | 52 Hu | $52 \mathrm{H}_{\text {L }}$ |
| GRID7 | $57 \mathrm{H}_{\mathrm{L}}$ | $56 \mathrm{H}_{\mathrm{U}}$ | $56 \mathrm{H}_{\mathrm{L}}$ | $55 \mathrm{H}_{\mathrm{U}}$ | $55 \mathrm{H}_{\mathrm{L}}$ |
| GRID8 | 5 AH | $59 \mathrm{H}_{\mathrm{U}}$ | 59 H | $58 \mathrm{H}_{\mathrm{U}}$ | $58 \mathrm{H}_{\mathrm{L}}$ |
| GRID9 | $5 \mathrm{DH}_{\mathrm{L}}$ | $5 \mathrm{CH}_{\mathrm{L}}$ | $5 \mathrm{CH}_{\mathrm{L}}$ | 5 BH L | $5 \mathrm{BH} \mathrm{L}^{\text {}}$ |
| GRID10 | $60 \mathrm{H}_{\mathrm{L}}$ | 5 FH L | 5 FH L | 5 EH L | $5 \mathrm{EH}_{\mathrm{L}}$ |
| GRID11 | $63 \mathrm{H}_{\mathrm{L}}$ | $62 \mathrm{H}_{\mathrm{L}}$ | $62 \mathrm{H}_{\mathrm{L}}$ | $61 \mathrm{H}_{\mathrm{L}}$ | $61 \mathrm{H}_{\text {L }}$ |
| GRID12 | $66 \mathrm{H}_{\mathrm{L}}$ | $65 \mathrm{H}_{\mathrm{L}}$ | $65 \mathrm{H}_{\mathrm{L}}$ | $64 \mathrm{H}_{\mathrm{L}}$ | $64 \mathrm{H}_{\mathrm{L}}$ |
| GRID13 | 69 H | $68 \mathrm{H}_{\mathrm{L}}$ | $68 \mathrm{H}_{\mathrm{L}}$ | $67 \mathrm{H}_{\mathrm{L}}$ | $67 \mathrm{H}_{\mathrm{L}}$ |
| GRID14 | $6 \mathrm{CH}_{\mathrm{L}}$ | 6 BH L | 6 BH L | $6 \mathrm{AH}_{\mathrm{L}}$ | $6 \mathrm{AH}_{\mathrm{L}}$ |
| GRID15 | 6 FH L | $6 \mathrm{EH}_{\mathrm{L}}$ | $6 \mathrm{EH}_{\mathrm{L}}$ | 6 DH L | $6 \mathrm{DH}_{L}$ |

Note: Only the lower 4 bits of the addresses assigned to SEG16 through SEG19 are valid, the higher 4 bits are ignored.

VFD Display Control Register - VFDC

| Bit No. | Label | Function |
| :---: | :---: | :---: |
| 3~0 | $\begin{aligned} & \text { VGS3~ } \\ & \text { VGS0 } \end{aligned}$ | Selects the VFD display mode $0 x x x=8$ grids, 20 segments 1000=9 grids, 19 segments 1001=10 grids, 18 segments 1010=11 grids, 17 segments 1011=12 grids, 16 segments 1100=13 grids, 15 segments 1101=14 grids, 14 segments $1110=15$ grids, 13 segments 1111=16 grids, 12 segments |
| 4 |  | Controls the VFD display (1=enable; 0=disable) |
| 7~5 | VDM2~ <br> VDMO | Sets the VFD dimming quantity $000=$ set pulse width to $1 / 16$. $001=$ set pulse width to $2 / 16$. $010=$ set pulse width to $4 / 16$. 011=set pulse width to 10/16. $100=$ set pulse width to $11 / 16$. $101=$ set pulse width to $12 / 16$. $110=$ set pulse width to $13 / 16$. $111=$ set pulse width to $14 / 16$. |

VFDC (28H) Register

At power-on, the 16 -grid, 12 -segment \& $1 / 16$ pulse width are set and the VFD display is disabled.

VFD clock source may come from the RTC or the system clock/4 (fsYS $/ 4$ ), if the $\mathrm{f}_{\mathrm{SYS}} / 4$ is selected to be the VFD clock source, when the system is in the HALT mode, the H/W should automatically turn the VFD off.

## Low Voltage Reset

There is a low voltage reset circuit (LVR) implemented in the microcontroller. This function can be enabled/disabled by options.

The LVR has the same effect or function with the external RES signal which performs a chip reset. During the HALT state, the LVR is disabled.

## Options

The following shows the configuration options in the HT49RV9/HT49CV9. All these options should be defined in order to ensure a properly functioning system.


| LVR selection. LVR has enable or disable option |
| :--- |
| PFD selection. If PA3 is set as a PFD output, there are two types of selections; One is PFD0 as the PFD output, the |
| other is PFD1 as the PFD output. PFD0, PFD1 are the timer overflow signals of the Timer/Event Counter 0, |
| Timer/Event Counter 1, respectively. |
| PWM selection: (7+1) or (6+2) mode |
| PD0: level output or PWM0 output |
| PD1: level output or PWM1 output |
| PD2: level output or PWM2 output |
| PD3: level output or PWM3 output |
| INT0 or $\overline{\text { INT1 }}$ triggering edge selection: disable; high to low; low to high; low to high or high to low. |
| SIOCLK: Serial interface clock. There are falling edge or rising edge |
| CSEN: Serial bus selection: enable or disable |
| WCOL: SBDR write conflict |

## Application Circuits



Note: The resistance and capacitance for reset circuit should be designed to ensure that the VDD is stable and remains in a valid range of the operating voltage before bringing $\overline{R E S}$ high.
"*" Make the length of the wiring, which is connected to the $\overline{R E S}$ pin as short as possible, to avoid noise interference.

The following table shows the C1, C2 and R1 values corresponding to the different crystal values. (For reference only)

| Crystal or Resonator | C1, C2 | R1 |
| :--- | :---: | :---: |
| 4 MHz Crystal | 0 pF | $10 \mathrm{k} \Omega$ |
| 4 MHz Resonator | 10 pF | $12 \mathrm{k} \Omega$ |
| 3.58 MHz Crystal | 0 pF | $10 \mathrm{k} \Omega$ |
| 3.58 MHz Resonator | 25 pF | $10 \mathrm{k} \Omega$ |
| 2 MHz Crystal \& Resonator | 25 pF | $10 \mathrm{k} \Omega$ |
| 1 MHz Crystal | 35 pF | $27 \mathrm{k} \Omega$ |
| 480 kHz Resonator | 300 pF | $9.1 \mathrm{k} \Omega$ |
| 455 kHz Resonator | 300 pF | $10 \mathrm{k} \Omega$ |
| 429 kHz Resonator | 300 pF | $10 \mathrm{k} \Omega$ |

The function of the resistor R1 is to ensure that the oscillator will switch off should low voltage conditions occur. Such a low voltage, as mentioned here, is one which is less than the lowest value of the MCU operating voltage. Note however that if the LVR is enabled then R1 can be removed.

Instruction Set Summary

| Mnemonic | Description | Instruction Cycle | Flag Affected |
| :---: | :---: | :---: | :---: |
| Arithmetic |  |  |  |
| ADD A,[m] | Add data memory to ACC | 1 | Z,C,AC,Ov |
| ADDM A,[m] | Add ACC to data memory | $1^{(1)}$ | Z,C,AC,OV |
| ADD A, x | Add immediate data to ACC | 1 | Z,C,AC,OV |
| ADC A,[m] | Add data memory to ACC with carry | 1 | Z,C,AC, ov |
| ADCM A,[m] | Add ACC to data memory with carry | $1^{(1)}$ | Z,C,AC,OV |
| SUB A, x | Subtract immediate data from ACC | 1 | Z,C,AC,OV |
| SUB A,[m] | Subtract data memory from ACC | 1 | Z,C,AC,OV |
| SUBM A,[m] | Subtract data memory from ACC with result in data memory | $1^{(1)}$ | Z,C,AC,OV |
| SBC A,[m] | Subtract data memory from ACC with carry | 1 | Z,C,AC,OV |
| SBCM A,[m] | Subtract data memory from ACC with carry and result in data memory | $1^{(1)}$ | Z,C,AC, ov |
| DAA [m] | Decimal adjust ACC for addition with result in data memory | $1^{(1)}$ | C |
| Logic Operation |  |  |  |
| AND A,[m] | AND data memory to ACC | 1 | Z |
| OR A,[m] | OR data memory to ACC | 1 | Z |
| XOR A,[m] | Exclusive-OR data memory to ACC | 1 | z |
| ANDM A,[m] | AND ACC to data memory | $1^{(1)}$ | Z |
| ORM A,[m] | OR ACC to data memory | $1^{(1)}$ | z |
| XORM A,[m] | Exclusive-OR ACC to data memory | $1^{(1)}$ | Z |
| AND A, x | AND immediate data to ACC | 1 | Z |
| OR A, x | OR immediate data to ACC | 1 | z |
| XOR A, x | Exclusive-OR immediate data to ACC | 1 | Z |
| CPL [m] | Complement data memory | $1^{(1)}$ | z |
| CPLA [m] | Complement data memory with result in ACC | 1 | z |
| Increment \& Decrement |  |  |  |
| INCA [m] | Increment data memory with result in ACC | 1 | Z |
| INC [m] | Increment data memory | $1^{(1)}$ | Z |
| DECA [m] | Decrement data memory with result in ACC | 1 | z |
| DEC [m] | Decrement data memory | $1^{(1)}$ | Z |
| Rotate |  |  |  |
| RRA [m] | Rotate data memory right with result in ACC | 1 | None |
| RR [m] | Rotate data memory right | $1^{(1)}$ | None |
| RRCA [m] | Rotate data memory right through carry with result in ACC | 1 | C |
| RRC [m] | Rotate data memory right through carry | $1^{(1)}$ | C |
| RLA [m] | Rotate data memory left with result in ACC | - | None |
| RL [m] | Rotate data memory left | $1^{(1)}$ | None |
| RLCA [m] | Rotate data memory left through carry with result in ACC | 1 | C |
| RLC [m] | Rotate data memory left through carry | $1^{(1)}$ | C |
| Data Move |  |  |  |
| MOV A,[m] | Move data memory to ACC |  |  |
| MOV [m],A | Move ACC to data memory | $1^{(1)}$ | None |
| MOV A, x | Move immediate data to ACC | 1 | None |
| Bit Operation |  |  |  |
| CLR [m].i | Clear bit of data memory | $1^{(1)}$ | None |
| SET [m].i | Set bit of data memory | $1^{(1)}$ | None |


| Mnemonic | Description | Instruction Cycle | Flag <br> Affected |
| :---: | :---: | :---: | :---: |
| Branch |  |  |  |
| ```JMP addr SZ [m] SZA [m] SZ [m].i SNZ [m].i SIZ [m] SDZ [m] SIZA [m] SDZA [m] CALL addr RET RET A,x RETI``` | Jump unconditionally <br> Skip if data memory is zero <br> Skip if data memory is zero with data movement to ACC <br> Skip if bit i of data memory is zero <br> Skip if bit i of data memory is not zero <br> Skip if increment data memory is zero <br> Skip if decrement data memory is zero <br> Skip if increment data memory is zero with result in ACC <br> Skip if decrement data memory is zero with result in ACC <br> Subroutine call <br> Return from subroutine <br> Return from subroutine and load immediate data to ACC <br> Return from interrupt | $\begin{gathered} 2 \\ 1^{(2)} \\ 1^{(2)} \\ 1^{(2)} \\ 1^{(2)} \\ 1^{(3)} \\ 1^{(3)} \\ 1^{(2)} \\ 1^{(2)} \\ 2 \\ 2 \\ 2 \\ 2 \end{gathered}$ | None None None None None None None None None None None None None |
| Table Read |  |  |  |
| TABRDC [m] TABRDL [m] | Read ROM code (current page) to data memory and TBLH Read ROM code (last page) to data memory and TBLH | $\begin{aligned} & 2^{(1)} \\ & 2^{(1)} \end{aligned}$ | None <br> None |
| Miscellaneous |  |  |  |
| NOP <br> CLR [m] <br> SET [m] <br> CLR WDT <br> CLR WDT1 <br> CLR WDT2 <br> SWAP [m] <br> SWAPA [m] <br> HALT | No operation <br> Clear data memory <br> Set data memory <br> Clear Watchdog Timer <br> Pre-clear Watchdog Timer <br> Pre-clear Watchdog Timer <br> Swap nibbles of data memory <br> Swap nibbles of data memory with result in ACC <br> Enter power down mode | $\begin{gathered} 1 \\ 1^{(1)} \\ 1^{(1)} \\ 1 \\ 1 \\ 1 \\ 1^{(1)} \\ 1 \\ 1 \end{gathered}$ | None None None TO,PDF TO $^{(4)}$, PDF $^{(4)}$ TO $^{(4)}$, PDF $^{(4)}$ None None TO,PDF |

Note: x: Immediate data
m : Data memory address
A: Accumulator
i: 0~7 number of bits
addr: Program memory address
$\checkmark$ : Flag is affected
-: Flag is not affected
${ }^{(1)}$ : If a loading to the PCL register occurs, the execution cycle of instructions will be delayed for one more cycle (four system clocks).
(2). If a skipping to the next instruction occurs, the execution cycle of instructions will be delayed for one more cycle (four system clocks). Otherwise the original instruction cycle is unchanged.
(3): (1) and ${ }^{(2)}$
${ }^{(4)}$ : The flags may be affected by the execution status. If the Watchdog Timer is cleared by executing the CLR WDT1 or CLR WDT2 instruction, the TO and PDF are cleared. Otherwise the TO and PDF flags remain unchanged.

## Instruction Definition

ADC A,[m]
Description

Operation
Affected flag(s)

## ADCM A,[m]

Description

Operation
Affected flag(s)

## ADD A,[m]

Description

Operation
Affected flag(s)

## ADD A,x

Description

Operation
Affected flag(s)

ADDM A,[m]
Description

Operation
Affected flag(s)

Add data memory and carry to the accumulator
The contents of the specified data memory, accumulator and the carry flag are added simultaneously, leaving the result in the accumulator.
$A C C \leftarrow A C C+[m]+C$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\sqrt{n}$ | $\sqrt{ }$ | $\sqrt{ }$ | $\sqrt{ }$ |

Add the accumulator and carry to data memory
The contents of the specified data memory, accumulator and the carry flag are added simultaneously, leaving the result in the specified data memory.
$[\mathrm{m}] \leftarrow \mathrm{ACC}+[\mathrm{m}]+\mathrm{C}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\sqrt{ }$ | $\checkmark$ | $\sqrt{ }$ |

Add data memory to the accumulator
The contents of the specified data memory and the accumulator are added. The result is stored in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}+[\mathrm{m}]$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\sqrt{ }$ |

Add immediate data to the accumulator
The contents of the accumulator and the specified data are added, leaving the result in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}+\mathrm{x}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ |

Add the accumulator to the data memory
The contents of the specified data memory and the accumulator are added. The result is stored in the data memory.
$[\mathrm{m}] \leftarrow \mathrm{ACC}+[\mathrm{m}]$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\sqrt{ }$ | $\checkmark$ | $\sqrt{ }$ |

AND A,[m]
Description

Operation
Affected flag(s)

## AND A,x

Description

Operation
Affected flag(s)

## ANDM A,[m]

Description

Operation
Affected flag(s)

CALL addr
Description

Operation

Affected flag(s)

## CLR [m]

Description
Operation
Affected flag(s)

Logical AND accumulator with data memory
Data in the accumulator and the specified data memory perform a bitwise logical_AND operation. The result is stored in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}$ "AND" $[\mathrm{m}]$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | $\sqrt{n}$ | - | - |

Logical AND immediate data to the accumulator
Data in the accumulator and the specified data perform a bitwise logical_AND operation. The result is stored in the accumulator.

ACC $\leftarrow$ ACC "AND" $x$

| TO | PDF | OV | Z |  | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | $\sqrt{2}$ | - | - |

Logical AND data memory with the accumulator
Data in the specified data memory and the accumulator perform a bitwise logical_AND operation. The result is stored in the data memory.
$[\mathrm{m}] \leftarrow \mathrm{ACC}$ "AND" $[\mathrm{m}]$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | $\sqrt{n}$ | - | - |

Subroutine call
The instruction unconditionally calls a subroutine located at the indicated address. The program counter increments once to obtain the address of the next instruction, and pushes this onto the stack. The indicated address is then loaded. Program execution continues with the instruction at this address.

Stack $\leftarrow$ Program Counter+1
Program Counter $\leftarrow$ addr

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |

Clear data memory
The contents of the specified data memory are cleared to 0 .
$[\mathrm{m}] \leftarrow \mathrm{OOH}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

CLR [m].i
Description
Operation
Affected flag(s)

CLR WDT
Description

Operation

Affected flag(s)

CLR WDT1
Description

Operation

Affected flag(s)

## CLR WDT2

Description

Operation

Affected flag(s)

CPL [m]
Description

Operation
Affected flag(s)

Clear bit of data memory
The bit i of the specified data memory is cleared to 0 .
$[\mathrm{m}] . \mathrm{i} \leftarrow 0$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | - | - | - |

Clear Watchdog Timer
The WDT is cleared (clears the WDT). The power down bit (PDF) and time-out bit (TO) are cleared.

WDT $\leftarrow 00 \mathrm{H}$
PDF and $\mathrm{TO} \leftarrow 0$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | - | - | - | - |

Preclear Watchdog Timer
Together with CLR WDT2, clears the WDT. PDF and TO are also cleared. Only execution of this instruction without the other preclear instruction just sets the indicated flag which implies this instruction has been executed and the TO and PDF flags remain unchanged.

WDT $\leftarrow 00 \mathrm{H}^{*}$
PDF and $\mathrm{TO} \leftarrow 0^{*}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $0^{*}$ | $0^{*}$ | - | - | - | - |

Preclear Watchdog Timer
Together with CLR WDT1, clears the WDT. PDF and TO are also cleared. Only execution of this instruction without the other preclear instruction, sets the indicated flag which implies this instruction has been executed and the TO and PDF flags remain unchanged.

WDT $\leftarrow 00 \mathrm{H}^{*}$
PDF and $\mathrm{TO} \leftarrow 0^{\text {* }}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $0^{*}$ | $0^{*}$ | - | - | - | - |

Complement data memory
Each bit of the specified data memory is logically complemented (1's complement). Bits which previously contained a 1 are changed to 0 and vice-versa.
$[\mathrm{m}] \leftarrow[\overline{\mathrm{m}}]$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | $\sqrt{c \mid}$ | - | - |



## HALT

Description

Operation

Affected flag(s)

INC [m]
Description
Operation
Affected flag(s)

INCA [m]
Description

Operation
Affected flag(s)

## JMP addr

Description

Operation
Affected flag(s)

MOV A,[m]
Description
Operation
Affected flag(s)

Directly jump
The program counter are replaced with the directly-specified address unconditionally, and control is passed to this destination.
Program Counter $\leftarrow$ addr

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Move data memory to the accumulator
The contents of the specified data memory are copied to the accumulator.
$\mathrm{ACC} \leftarrow[\mathrm{m}]$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

MOV A, $x$
Description
Operation
Affected flag(s)

MOV [m],A
Description

Operation
Affected flag(s)

NOP
Description
Operation
Affected flag(s)

OR A,[m]
Description

Operation
Affected flag(s)

OR A, $x$
Description

Operation
Affected flag(s)

ORM A,[m]
Description

Operation
Affected flag(s)

Move immediate data to the accumulator
The 8-bit data specified by the code is loaded into the accumulator.
ACC $\leftarrow \mathrm{x}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Move the accumulator to data memory
The contents of the accumulator are copied to the specified data memory (one of the data memories).
[m] $\leftarrow$ ACC

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

No operation
No operation is performed. Execution continues with the next instruction.
Program Counter $\leftarrow$ Program Counter+1

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Logical OR accumulator with data memory
Data in the accumulator and the specified data memory (one of the data memories) perform a bitwise logical_OR operation. The result is stored in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}$ " OR " $[\mathrm{m}]$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | $\sqrt{n}$ | - | - |

Logical OR immediate data to the accumulator
Data in the accumulator and the specified data perform a bitwise logical_OR operation. The result is stored in the accumulator.

ACC $\leftarrow A C C$ "OR" $x$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | $\sqrt{n}$ | - | - |

Logical OR data memory with the accumulator
Data in the data memory (one of the data memories) and the accumulator perform a bitwise logical_OR operation. The result is stored in the data memory.
$[\mathrm{m}] \leftarrow \mathrm{ACC}$ "OR" [m]

| TO | PDF | OV | Z |  | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | $\sqrt{2}$ | - | - |

RET
Description
Operation
Affected flag(s)

RET A,x
Description

Operation

Affected flag(s)

RETI
Description

Operation

Affected flag(s)

RL [m]
Description
Operation

Affected flag(s)

RLA [m]
Description

Operation

Affected flag(s)

Return from subroutine
The program counter is restored from the stack. This is a 2-cycle instruction.
Program Counter $\leftarrow$ Stack

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Return and place immediate data in the accumulator
The program counter is restored from the stack and the accumulator loaded with the specified 8-bit immediate data.

Program Counter $\leftarrow$ Stack ACC $\leftarrow x$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Return from interrupt
The program counter is restored from the stack, and interrupts are enabled by setting the EMI bit. EMI is the enable master (global) interrupt bit.

Program Counter $\leftarrow$ Stack
EMI $\leftarrow 1$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Rotate data memory left
The contents of the specified data memory are rotated 1 bit left with bit 7 rotated into bit 0 .
$[\mathrm{m}] .(\mathrm{i}+1) \leftarrow[\mathrm{m}] . \mathrm{i} ;[\mathrm{m}] . \mathrm{i}$.bit i of the data memory ( $\mathrm{i}=0 \sim 6$ )
$[\mathrm{m}] .0 \leftarrow[\mathrm{~m}] .7$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Rotate data memory left and place result in the accumulator
Data in the specified data memory is rotated 1 bit left with bit 7 rotated into bit 0 , leaving the rotated result in the accumulator. The contents of the data memory remain unchanged.

ACC. $(\mathrm{i}+1) \leftarrow[\mathrm{m}] . \mathrm{i} ;[\mathrm{m}] . \mathrm{i}$ :bit i of the data memory ( $\mathrm{i}=0 \sim 6$ )
ACC. $0 \leftarrow[\mathrm{~m}] .7$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

RLC [m]
Description

Operation

Affected flag(s)

RLCA [m]
Description

Operation

Affected flag(s)

RR [m]
Description
Operation

Affected flag(s)

RRA [m]
Description

Operation

Affected flag(s)

## RRC [m]

Description

Operation

Affected flag(s)

Rotate data memory left through carry
The contents of the specified data memory and the carry flag are rotated 1 bit left. Bit 7 replaces the carry bit; the original carry flag is rotated into the bit 0 position.
$[\mathrm{m}] .(\mathrm{i}+1) \leftarrow[\mathrm{m}] . \mathrm{i} ;[\mathrm{m}] . \mathrm{i}:$ bit i of the data memory ( $\mathrm{i}=0 \sim 6$ )
$[\mathrm{m}] .0 \leftarrow \mathrm{C}$
$C \leftarrow[\mathrm{~m}] .7$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | $V$ |

Rotate left through carry and place result in the accumulator
Data in the specified data memory and the carry flag are rotated 1 bit left. Bit 7 replaces the carry bit and the original carry flag is rotated into bit 0 position. The rotated result is stored in the accumulator but the contents of the data memory remain unchanged.
ACC. $(\mathrm{i}+1) \leftarrow[\mathrm{m}] . \mathrm{i} ;[\mathrm{m}] . \mathrm{i}$.bit i of the data memory ( $\mathrm{i}=0 \sim 6$ )
ACC. $0 \leftarrow C$
$C \leftarrow[\mathrm{~m}] .7$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | $V$ |

Rotate data memory right
The contents of the specified data memory are rotated 1 bit right with bit 0 rotated to bit 7 .
$[\mathrm{m}] . \mathrm{i} \leftarrow[\mathrm{m}] .(\mathrm{i}+1) ;[\mathrm{m}] . \mathrm{i}$ ibit i of the data memory $(\mathrm{i}=0 \sim 6)$
$[\mathrm{m}] .7 \leftarrow[\mathrm{~m}] .0$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Rotate right and place result in the accumulator
Data in the specified data memory is rotated 1 bit right with bit 0 rotated into bit 7 , leaving the rotated result in the accumulator. The contents of the data memory remain unchanged.
ACC. (i) $\leftarrow[\mathrm{m}]$. (i+1); [m].i:bit i of the data memory ( $\mathrm{i}=0 \sim 6$ )
ACC. $7 \leftarrow[\mathrm{~m}] .0$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Rotate data memory right through carry
The contents of the specified data memory and the carry flag are together rotated 1 bit right. Bit 0 replaces the carry bit; the original carry flag is rotated into the bit 7 position.
$[\mathrm{m}] . \mathrm{i} \leftarrow[\mathrm{m}] .(\mathrm{i}+1) ;[\mathrm{m}] . \mathrm{i}$.bit i of the data memory $(\mathrm{i}=0 \sim 6)$
$[\mathrm{m}] .7 \leftarrow \mathrm{C}$
$C \leftarrow[\mathrm{~m}] .0$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | - | - | $V$ |

RRCA [m]
Description
Operation
Affected flag(s)

SBC A,[m]
Description

Operation
Affected flag(s)

SBCM A,[m]
Description

Operation
Affected flag(s)

SDZ [m]
Description

Operation
Affected flag(s)

## SDZA [m]

Description

Operation
Affected flag(s)

Rotate right through carry and place result in the accumulator
Data of the specified data memory and the carry flag are rotated 1 bit right. Bit 0 replaces the carry bit and the original carry flag is rotated into the bit 7 position. The rotated result is stored in the accumulator. The contents of the data memory remain unchanged.
ACC. $\mathrm{i} \leftarrow[\mathrm{m}] .(\mathrm{i}+1)$; [m]. i .bit i of the data memory ( $\mathrm{i}=0 \sim 6$ )
ACC. $7 \leftarrow \mathrm{C}$
$C \leftarrow[\mathrm{~m}] .0$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | $\checkmark$ |

Subtract data memory and carry from the accumulator
The contents of the specified data memory and the complement of the carry flag are subtracted from the accumulator, leaving the result in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}+[\overline{\mathrm{m}}]+\mathrm{C}$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\sqrt{ }$ |

Subtract data memory and carry from the accumulator
The contents of the specified data memory and the complement of the carry flag are subtracted from the accumulator, leaving the result in the data memory.

$$
[\mathrm{m}] \leftarrow \mathrm{ACC}+[\overline{\mathrm{m}}]+\mathrm{C}
$$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\checkmark$ |

Skip if decrement data memory is 0
The contents of the specified data memory are decremented by 1 . If the result is 0 , the next instruction is skipped. If the result is 0 , the following instruction, fetched during the current instruction execution, is discarded and a dummy cycle is replaced to get the proper instruction ( 2 cycles). Otherwise proceed with the next instruction (1 cycle).
Skip if $([m]-1)=0,[m] \leftarrow([m]-1)$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Decrement data memory and place result in ACC, skip if 0
The contents of the specified data memory are decremented by 1 . If the result is 0 , the next instruction is skipped. The result is stored in the accumulator but the data memory remains unchanged. If the result is 0 , the following instruction, fetched during the current instruction execution, is discarded and a dummy cycle is replaced to get the proper instruction (2 cycles). Otherwise proceed with the next instruction (1 cycle).

Skip if $([m]-1)=0, A C C \leftarrow([m]-1)$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |



SUB A,[m]
Description

Operation
Affected flag(s)

SUBM A,[m]
Description

Operation
Affected flag(s)

## SUB A,x

Description

Operation
Affected flag(s)

SWAP [m]
Description

Operation
Affected flag(s)

SWAPA [m]
Description

Operation

Affected flag(s)

Subtract data memory from the accumulator
The specified data memory is subtracted from the contents of the accumulator, leaving the result in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}+[\overline{\mathrm{m}}]+1$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | $\sqrt{n}$ | $\sqrt{ }$ | $\sqrt{ }$ | $\sqrt{ }$ |

Subtract data memory from the accumulator
The specified data memory is subtracted from the contents of the accumulator, leaving the result in the data memory.
$[\mathrm{m}] \leftarrow \mathrm{ACC}+[\overline{\mathrm{m}}]+1$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\sqrt{ }$ |

Subtract immediate data from the accumulator
The immediate data specified by the code is subtracted from the contents of the accumulator, leaving the result in the accumulator.
$A C C \leftarrow A C C+\bar{x}+1$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | $\checkmark$ | $\checkmark$ | $\checkmark$ | $\sqrt{ }$ |

Swap nibbles within the data memory
The low-order and high-order nibbles of the specified data memory (1 of the data memories) are interchanged.
$[\mathrm{m}] .3 \sim[\mathrm{~m}] .0 \leftrightarrow[\mathrm{~m}] .7 \sim[\mathrm{~m}] .4$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | - | - | - |

Swap data memory and place result in the accumulator
The low-order and high-order nibbles of the specified data memory are interchanged, writing the result to the accumulator. The contents of the data memory remain unchanged.

ACC. $3 \sim$ ACC. $0 \leftarrow[\mathrm{~m}] .7 \sim[\mathrm{~m}] .4$
ACC. $7 \sim$ ACC. $4 \leftarrow[\mathrm{~m}] .3 \sim[\mathrm{~m}] .0$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

SZ [m]

Operation
Affected flag(s)

## TABRDC [m]

Description

Operation

Affected flag(s)

TABRDL [m]
Description

Operation

Affected flag(s)


Skip if data memory is 0
If the contents of the specified data memory are 0 , the following instruction, fetched during the current instruction execution, is discarded and a dummy cycle is replaced to get the proper instruction (2 cycles). Otherwise proceed with the next instruction (1 cycle).
Skip if $[m]=0$

Move data memory to ACC, skip if 0
The contents of the specified data memory are copied to the accumulator. If the contents is 0 , the following instruction, fetched during the current instruction execution, is discarded and a dummy cycle is replaced to get the proper instruction (2 cycles). Otherwise proceed with the next instruction (1 cycle).

Skip if $[m]=0$

Skip if bit $i$ of the data memory is 0
If bit $i$ of the specified data memory is 0 , the following instruction, fetched during the current tion (2 cycles). Otherwise proceed with the next instruction (1 cycle).
Skip if [m]. $\mathrm{i}=0$

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Move the ROM code (current page) to TBLH and data memory
The low byte of ROM code (current page) addressed by the table pointer (TBLP) is moved to the specified data memory and the high byte transferred to TBLH directly.
[m] $\leftarrow$ ROM code (low byte) TBLH $\leftarrow$ ROM code (high byte)

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

Move the ROM code (last page) to TBLH and data memory
The low byte of ROM code (last page) addressed by the table pointer (TBLP) is moved to the data memory and the high byte transferred to TBLH directly.
[m] $\leftarrow$ ROM code (low byte) TBLH $\leftarrow$ ROM code (high byte)

| TO | PDF | OV | Z | AC | C |
| :---: | :---: | :---: | :---: | :---: | :---: |
| - | - | - | - | - | - |

XOR A,[m]
Description

Operation
Affected flag(s)

XORM A,[m]
Description

Operation
Affected flag(s)

## XOR A,x

Description

Operation
Affected flag(s)

Logical XOR accumulator with data memory
Data in the accumulator and the indicated data memory perform a bitwise logical Exclusive_OR operation and the result is stored in the accumulator.
$\mathrm{ACC} \leftarrow \mathrm{ACC}$ "XOR" $[\mathrm{m}]$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | $\sqrt{n}$ | - | - |

Logical XOR data memory with the accumulator
Data in the indicated data memory and the accumulator perform a bitwise logical Exclusive_OR operation. The result is stored in the data memory. The 0 flag is affected.
$[\mathrm{m}] \leftarrow \mathrm{ACC}$ "XOR" $[\mathrm{m}]$

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C |  |  |  |  |  |
| - | - | - | $\sqrt{c \mid}$ | - | - |

Logical XOR immediate data to the accumulator
Data in the accumulator and the specified data perform a bitwise logical Exclusive_OR operation. The result is stored in the accumulator. The 0 flag is affected.
ACC $\leftarrow A C C$ "XOR" x

| TO |  | PDF | OV | Z | AC |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | C |  |  |  |  |
| - | - | - | $\sqrt{2}$ | - | - |

## Package Information

100 -pin QFP ( $14 \times 20$ ) Outline Dimensions


| Symbol | Dimensions in mm |  |  |
| :---: | :---: | :---: | :---: |
|  | Min. | Nom. | Max. |
| A | 18.50 | - | 19.20 |
| B | 13.90 | - | 14.10 |
| C | 24.50 | - | 25.20 |
| D | 19.90 | - | 20.10 |
| E | - | 0.65 | - |
| F | - | 0.30 | - |
| G | 2.50 | - | 3.10 |
| H | - | - | 3.40 |
| I | - | 0.10 | - |
| J | 1 | - | 1.40 |
| K | 0.10 | - | 0.20 |
| $\alpha$ | $0^{\circ}$ | - | $7^{\circ}$ |

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